

Reduction

$L \leq_m L'$
We know L
is not in,
say, RE

this is the language
under investigation

L' turns out to be
outside of RE,
because of the
reduction

mistakes:

$L' \leq_m L$
not in RE
?? X Don't do this !!

A reduction is a sort of algorithm

$$L \leq_m L'$$

transform M
into M'

input: instance of L

enc(M)

output: instance of L'

enc(M')

many/many relations:

→ if $enc(M) \in L$ then $enc(M') \in L'$

if $enc(M) \notin L$ then $enc(M') \notin L'$

Let M_1, M_2 be TM. We write $\underline{\text{enc}(M_1, M_2)}$ =
encoding of the two TMs.

$$L = \left\{ \text{enc}(M_1, M_2) \mid \overline{L(M_1)} = L(M_2) \right\}$$

← string

Show that L is not in RE.

$$L_d, \underline{L_e} \quad L_e \leq_m L$$

Therefore, for some n , L is not in RE.

specify the reduction

input: $enc(M)$

→ output: $enc(M_1, M_2)$

source lang.

$L \leq_m L$

target lang.

these TMs depend
on M !!

possible mistakes:

M_1, M_2 independent of M |
Don't do !!

We want: if $L(M) = \emptyset$ then $L(M_1) = L(M_2)$

this is the property defining the source lang.

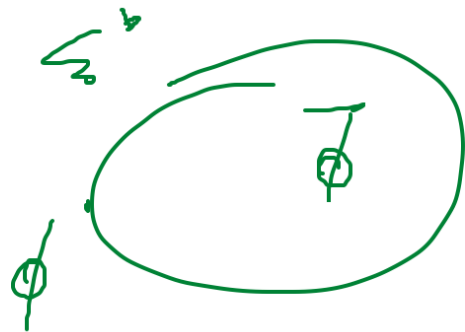
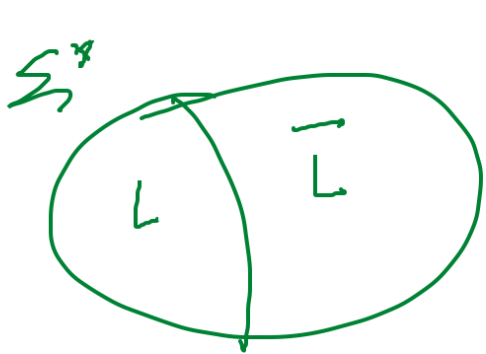
this is the property defining L_1 (target lang.)

Idea

set $M_1 = M$

if M recognizes empty language, $L(M) = \emptyset$

then $\overline{L(M_1)} = \Sigma^*$



$$\Sigma^* = \overline{\emptyset}$$

set $M_2 = M_{\Sigma^*}$ where $L(M_{\Sigma^*}) = \Sigma^*$

M_2 does not depend on M

but M_1 depends on M

input

output

M_1

copy

M



$$L(M_{\Sigma^*}) = \Sigma^*$$

$$\underline{M_2 = M_{\Sigma^*}}$$

(A) y^2y

if $L(M) = \emptyset$

r

this means

$$\underline{\text{enc}(M) \in L_e}$$

then

$$L(M_1) = \emptyset$$

then

$$\underline{L(M_1) = \Sigma^*}$$

then

$$\underline{L(M_1) = L(M_2)}$$

then

$$\text{enc}(M_1, M_2) \in L$$

(B) $n \geq n$

if $enc(M) \notin L_e$

then

$L(M) \neq \emptyset$

then

$L(M_1) \neq \emptyset$

then

$\overline{L(M_1)} \neq \Sigma^*$

then

$\overline{L(M_1)} \neq L(M_2)$

then

$enc(M_1, M_2) \notin L$

because M_2 recog.
 Σ^*

Σ^*



$L(M_1)$ $\overline{L(M_1)} \neq \Sigma^*$

We have shown $y \geq y \neq n \geq n$ property
Therefore the reduction is valid