Master Degree in Computer Engineering

## Final Exam for Automata, Languages and Computation

July 4th, 2024

- 1. [4 points] Consider the regular expression  $R = (ab + ba)^* \emptyset(aa)$ . Convert R into an equivalent  $\epsilon$ -NFA using the construction provided in the textbook, and report all the intermediate steps. Important: do not simplify the regular expression R before applying the construction.
- 2. [9 points] Let  $\Sigma = \{a, b, c\}$ . For  $w \in \Sigma^*$  and  $X \in \Sigma$ , we write  $\#_X(w)$  to denote the number of occurrences of X in w. Consider the following languages

$$L_1 = \{ w \mid w \in \Sigma^*, \ \#_a(w) = \#_b(w) = \#_c(w) \} ;$$
  

$$L_2 = \{ w \mid w \in \Sigma^*, \ \#_a(w) = \#_c(w) \} .$$

- (a) Prove that  $L_1$  is outside of CFL.
- (b) Prove that  $L_2$  is in CFL.
- (c) Prove that  $L_2$  is not in REG.
- 3. [5 points] Consider the CFG G implicitly defined by the following productions:

$$\begin{array}{l} S \rightarrow AAB \mid ABB \mid BBB \\ A \rightarrow aAB \mid bBB \\ B \rightarrow b \mid \varepsilon \end{array}$$

Perform on G the transformations indicated below, that have been specified in the textbook, in the given order. Report the CFGs obtained at each of the intermediate steps.

- (a) Eliminate the  $\varepsilon$ -productions
- (b) Eliminate the unary productions
- (c) Eliminate the useless symbols
- (d) Produce a CFG in Chomsky normal form equivalent to G.

(please turn to the next page)

- 4. **[5 points]** Assess whether the following statements are true or false, providing motivations for all of your answers.
  - (a) Let  $L_1$  be a language in REG with  $L_1$  non-finite, and let  $L_2$  be a language in CFL\REG. The language  $L_1 \cap L_2$  may be in CFL\REG.
  - (b) Let  $L_1$  be a language in REG with  $L_1$  non-finite, and let  $L_2$  be a language in CFL\REG. The language  $L_1 \cap L_2$  may be in REG.
  - (c) Let  $L_1, L_2$  be languages in CFL. The language  $L_1 \cap L_2$  belongs to  $\mathcal{P}$ , the class of languages that can be recognized in polynomial time by a TM.
  - (d) Let R be the string reversal operator, which we extend to languages. Let L be a language in REC. Then  $L^R$  belongs to REC.
- 5. [4 points] Define the diagonalization language  $L_d$ . Show that  $L_d$  is not an RE language, using the proof reported in the textbook.
- 6. [6 points] Consider the following property of the RE languages defined over the alphabet  $\Sigma = \{0, 1\}$ :

 $\mathcal{P} = \{L \mid L \in \text{RE, every string in } L \text{ has even length} \}$ 

and define  $L_{\mathcal{P}} = \{\mathsf{enc}(M) \mid L(M) \in \mathcal{P}\}.$ 

- (a) Use Rice's theorem to show that  $L_{\mathcal{P}}$  is not in REC.
- (b) Prove that  $L_{\mathcal{P}}$  is not in RE.