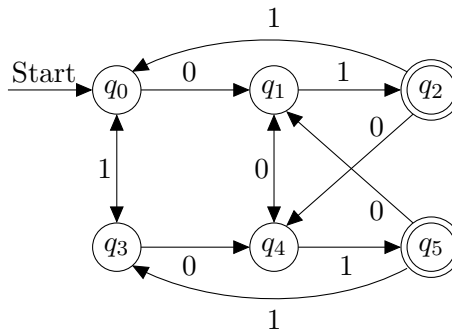


**Final Exam for
Automata, Languages and Computation**

September 13th, 2023

1. [6 points] Consider the DFA A whose transition function is graphically represented as follows (arcs with double direction represent two arcs in opposite directions)



- (a) Apply to A the tabular algorithm presented in the textbook for detecting pairs of equivalent states, reporting all the **intermediate steps**.
- (b) Specify the minimal DFA equivalent to A .
2. [7 points] Consider the following languages, defined over the alphabet $\Sigma = \{a, b\}$:

$$L_1 = \{a^m b a^n b a^p \mid m, n, p \geq 1, m < n < p\}$$

$$L_2 = \{a^m b a^n b a^p \mid m, n, p \geq 1, m + n < n + p\}$$

For each of the above languages, state whether it belongs to the class CFL and provide a mathematical proof of your answer.

3. [5 points] With reference to the membership problem for context-free languages, answer the following two questions.
- (a) Specify the dynamic programming algorithm developed in the textbook for the solution of this problem.
- (b) Consider the CFG G defined by the following rules:

$$S \rightarrow BD$$

$$B \rightarrow BB \mid b$$

$$D \rightarrow DD \mid d$$

Assuming as input the CFG G and the string $w = bbbdd$, trace the application of the above algorithm.

(please see next page)

4. **[6 points]** Assess whether the following statements are true or false, providing motivations for all of your answers.
- (a) The concatenation of a regular language and a context-free language is never a regular language.
 - (b) The concatenation of a regular language and a context-free language is always a regular language.
 - (c) The concatenation of a regular language and a context-free language is always a context-free language.
 - (d) The concatenation of a regular language and a language in \mathcal{P} is always a language in \mathcal{P} (\mathcal{P} is the class of languages that can be recognized in polynomial time by a TM).
5. **[4 points]** Define the diagonalization language L_d . Show that L_d is not an RE language using the proof reported in the textbook.
6. **[6 points]** Recall that for two strings $x, w \in \Sigma^*$, we say that x is an **infix** of w if we can write $w = uxv$ for some strings $u, v \in \Sigma^*$. Consider the following property of the RE languages defined over the alphabet $\Sigma = \{0, 1\}$:

$$\mathcal{P} = \{L \mid L \in \text{RE}, \text{ every string in } L \text{ has an infix } 01110\}$$

and define $L_{\mathcal{P}} = \{\text{enc}(M) \mid L(M) \in \mathcal{P}\}$.

- (a) Use Rice's theorem to show that $L_{\mathcal{P}}$ is not in REC.
- (b) Prove that $L_{\mathcal{P}}$ is not in RE.