

Probabilistic Model Checking

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# Lecture 3

## Discrete-time Markov Chains

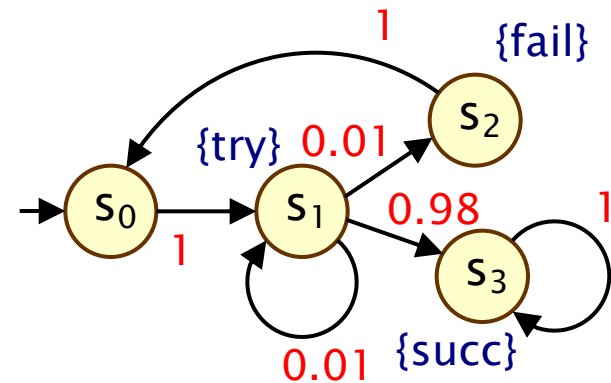
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# Discrete-time Markov chains

- State-transition systems augmented with probabilities
- States
  - **set of states** representing possible configurations of the system being modelled
- Transitions
  - transitions between states model evolution of system's state; occur in **discrete time-steps**
- Probabilities
  - probabilities of making transitions between states are given by **discrete probability distributions**



# Overview

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- Previous lecture: path-based properties, probabilistic reachability
- Transient state probabilities
- Long-run / steady-state probabilities
- Qualitative properties
  - repeated reachability
  - persistence

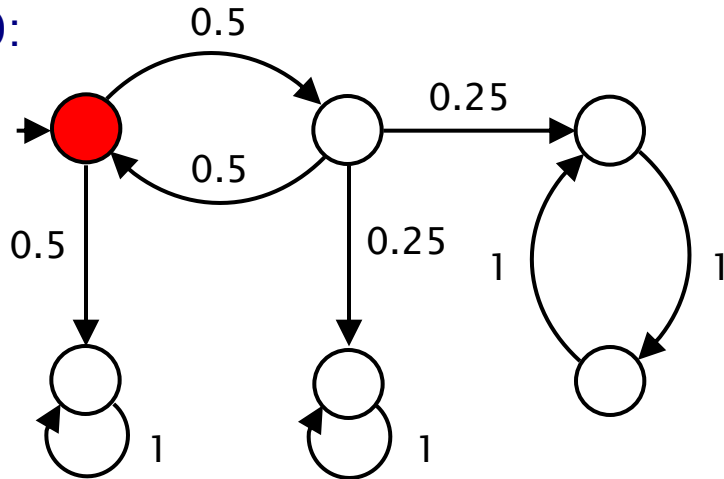
# Transient state probabilities

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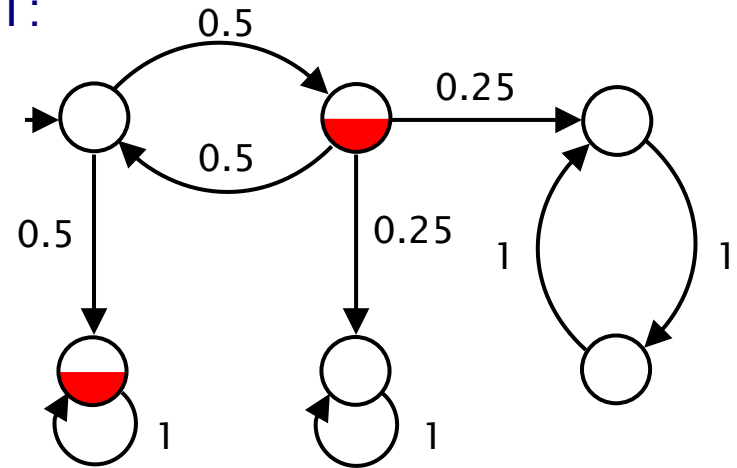
- What is the probability, having started in state  $s$ , of being in state  $s'$  at time  $k$ ?
  - i.e. after exactly  $k$  steps/transitions have occurred
  - this is the **transient state probability**:  $\pi_{s,k}(s')$
- **Transient state distribution**:  $\underline{\pi}_{s,k}$ 
  - (row) vector  $\underline{\pi}_{s,k}$  i.e.  $\pi_{s,k}(s')$  for all states  $s'$
- **Note**: this is a **discrete** probability distribution
  - so we have  $\underline{\pi}_{s,k} : S \rightarrow [0,1]$
  - recall instead  $\Pr_s : \Sigma_{\text{Path}(s)} \rightarrow [0,1]$ , where  $\Sigma_{\text{Path}(s)} \subseteq 2^{\text{Path}(s)}$

# Transient distributions

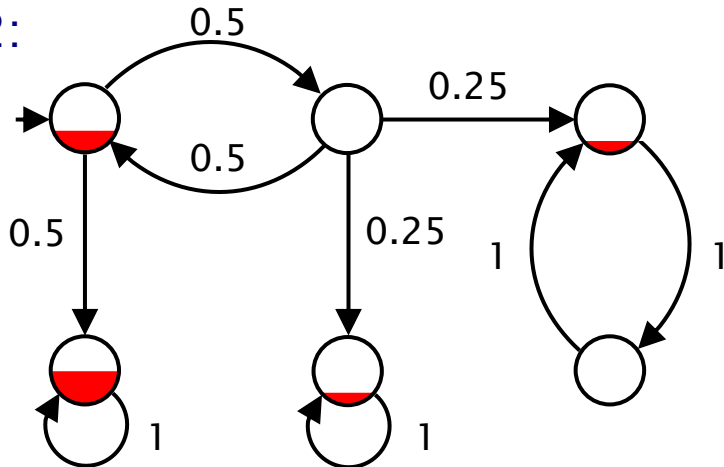
k=0:



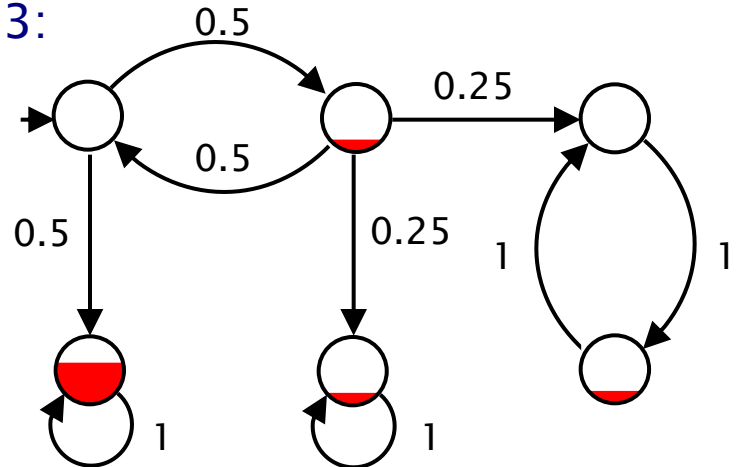
k=1:



k=2:



k=3:

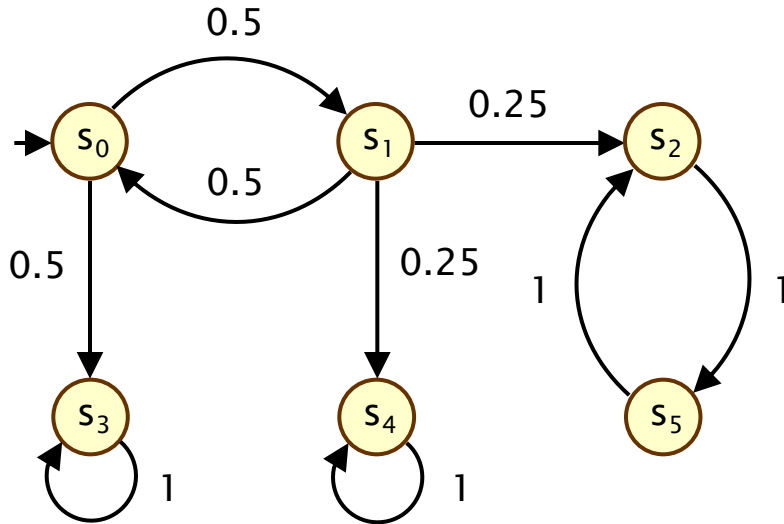


# Computing transient probabilities

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- Transient state probabilities:
  - $\pi_{s,k}(s') = \sum_{s'' \in S} \pi_{s,k-1}(s'') \cdot P(s'',s')$
  - (i.e. look at incoming transitions, into  $s'$ )
- Computation of transient state distribution:
  - $\underline{\pi}_{s,0}$  is the initial probability distribution
  - e.g. in our case  $\underline{\pi}_{s,0}(s') = 1$  if  $s'=s$  and  $\underline{\pi}_{s,0}(s') = 0$  otherwise
  - $\underline{\pi}_{s,k} = \underline{\pi}_{s,k-1} \cdot P$
- i.e. successive vector–matrix multiplications

# Computing transient probabilities



$$P = \begin{pmatrix} 0 & 0.5 & 0 & 0.5 & 0 & 0 \\ 0.5 & 0 & 0.25 & 0 & 0.25 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 \end{pmatrix}$$

$$\underline{\pi}_{s_0,0} = [1, 0, 0, 0, 0, 0]$$

$$\underline{\pi}_{s_0,1} = \left[0, \frac{1}{2}, 0, \frac{1}{2}, 0, 0\right]$$

$$\underline{\pi}_{s_0,2} = \left[\frac{1}{4}, 0, \frac{1}{8}, \frac{1}{2}, \frac{1}{8}, 0\right]$$

$$\underline{\pi}_{s_0,3} = \left[0, \frac{1}{8}, 0, \frac{5}{8}, \frac{1}{8}, \frac{1}{8}\right]$$

...

# Computing transient probabilities

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- $\underline{\pi}_{s,k} = \underline{\pi}_{s,k-1} \cdot \mathbf{P} = \underline{\pi}_{s,0} \cdot \mathbf{P}^k$
- $k^{\text{th}}$  matrix power:  $\mathbf{P}^k$ 
  - $\mathbf{P}$  gives one-step transition probabilities
  - $\mathbf{P}^k$  gives probabilities of  $k$ -step transition probabilities
  - i.e.  $\mathbf{P}^k(s,s') = \pi_{s,k}(s')$
- A possible optimisation: iterative squaring
  - e.g.  $\mathbf{P}^8 = ((\mathbf{P}^2)^2)^2$
  - only requires  $\log k$  multiplications
  - but potentially inefficient, e.g. if  $\mathbf{P}$  is large and sparse
  - in practice, successive vector-matrix multiplications preferred



# Notion of time in DTMCs

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- Two possible views on the timing aspects of a system modelled as a DTMC:
  1. Discrete time steps model time accurately
    - e.g. clock ticks in a model of an embedded device
    - or like dice example: interested in number of steps (tosses)
  2. Time–abstract model
    - no information assumed about the time transitions take
    - e.g. simple Zeroconf model
- In both cases, often beneficial to study long–run behaviour

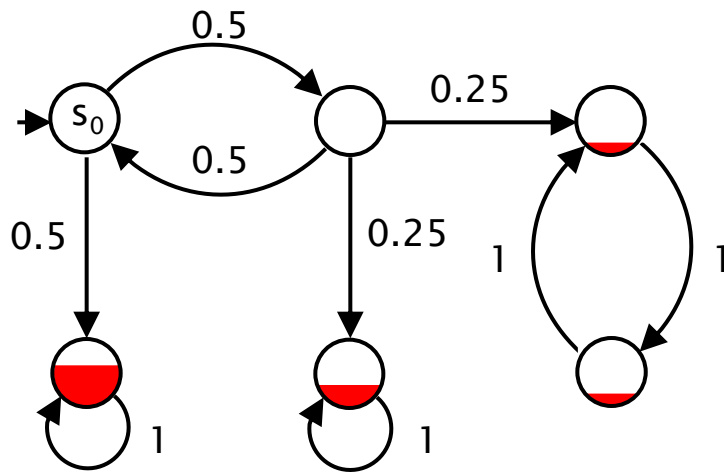
# Long-run behaviour

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- Consider the limit:  $\underline{\pi}_s = \lim_{k \rightarrow \infty} \underline{\pi}_{s,k}$ 
  - where  $\underline{\pi}_{s,k}$  is the transient state distribution at time  $k$ , having started in state  $s$
  - this limit, where it exists, is called the **limiting distribution**
  - steady-state of the model
- Intuitive idea
  - the percentage of time, in the long run, spent in each state
  - e.g. reliability: “in the long-run, what portion of time is the system in an operational state”

# Limiting distribution

- Example:



$$\underline{\pi}_{s_0,0} = [1,0,0,0,0,0]$$

$$\underline{\pi}_{s_0,1} = \left[0, \frac{1}{2}, 0, \frac{1}{2}, 0, 0\right]$$

$$\underline{\pi}_{s_0,2} = \left[\frac{1}{4}, 0, \frac{1}{8}, \frac{1}{2}, \frac{1}{8}, 0\right]$$

$$\underline{\pi}_{s_0,3} = \left[0, \frac{1}{8}, 0, \frac{5}{8}, \frac{1}{8}, \frac{1}{8}\right]$$

...

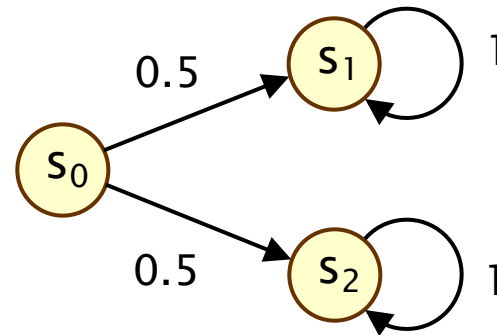
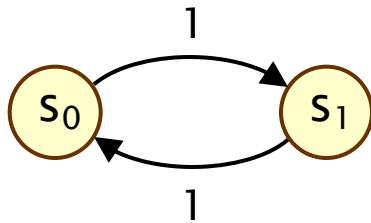
$$\underline{\pi}_{s_0} = \left[0, 0, \frac{1}{12}, \frac{2}{3}, \frac{1}{6}, \frac{1}{12}\right]$$

# Long-run behaviour

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- Questions:

- when does this limit exist?
- does it depend on the initial state/distribution?



- Need to consider underlying graph

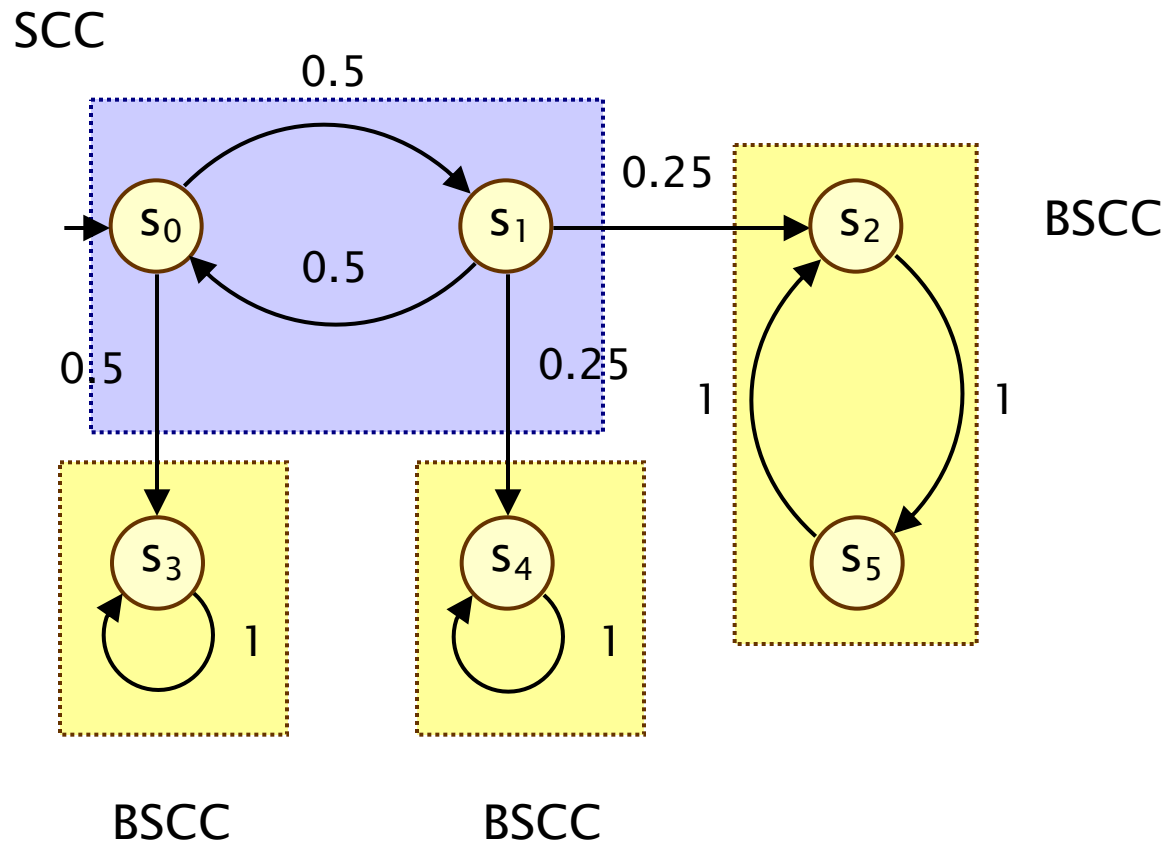
- $(V,E)$  where  $V$  are vertices and  $E \subseteq V \times V$  are edges
- $V = S$  and  $E = \{ (s,s') \text{ s.t. } P(s,s') > 0 \}$

# Graph terminology

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- A state  $s'$  is **reachable** from  $s$  if there is a finite path starting in  $s$  and ending in  $s'$
- A subset  $T$  of  $S$  is **strongly connected** if, for each pair of states  $s$  and  $s'$  in  $T$ ,  $s'$  is reachable from  $s$  passing only through states in  $T$
- A **strongly connected component** (SCC) is a *maximally* strongly connected set of states (i.e. no superset of it is also strongly connected)
- A **bottom strongly connected component** (BSCC) is an SCC  $T$  from which no state outside  $T$  is reachable from  $T$
- Alternative terminology: “ $s$  communicates with  $s'$ ”, “communicating class”, “recurrent class”

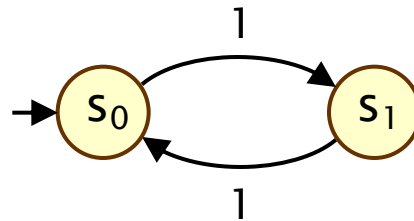
# Example – (B)SCCs



# Graph terminology

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- Markov chain is **irreducible** if all its states belong to a single BSCC; otherwise reducible



- A state  $s$  is **periodic**, with period  $d$ , if
  - the greatest common divisor of the set  $\{ n \mid f_s^{(n)} > 0 \}$  equals  $d$
  - where  $f_s^{(n)}$  is the probability of, when starting in state  $s$ , returning to state  $s$  in exactly  $n$  steps
- A Markov chain is **aperiodic** if all its states have period 1

# Steady-state probabilities

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- For a finite, irreducible, aperiodic DTMC (a.k.a., ergodic)
  - limiting distribution **always** exists
  - and is **independent** of initial state/distribution
- These are known as steady-state probabilities
  - (or equilibrium probabilities)
  - effect of initial distribution has disappeared, denoted  $\underline{\pi}$
- These probabilities can be computed as the unique solution of the linear equation system:

$$\underline{\pi} \cdot \mathbf{P} = \underline{\pi} \quad \text{and} \quad \sum_{s \in S} \underline{\pi}(s) = 1$$



# Steady-state – Balance equations

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- Known as **balance equations**

$$\underline{\pi} \cdot \mathbf{P} = \underline{\pi} \quad \text{and} \quad \sum_{s \in S} \underline{\pi}(s) = 1$$

- That is:

- $\underline{\pi}(s') = \sum_{s \in S} \underline{\pi}(s) \cdot \mathbf{P}(s, s')$

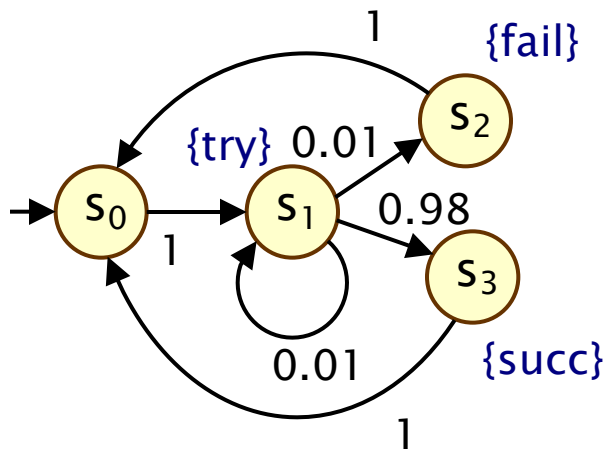
balance the probability of leaving and entering a state  $s'$

- $\sum_{s \in S} \underline{\pi}(s) = 1$

normalisation

# Steady-state – Example

- Let  $\underline{x} = \underline{\pi}$
- Solve:  $\underline{x} \cdot \mathbf{P} = \underline{x}$ ,  $\sum_s \underline{x}(s) = 1$



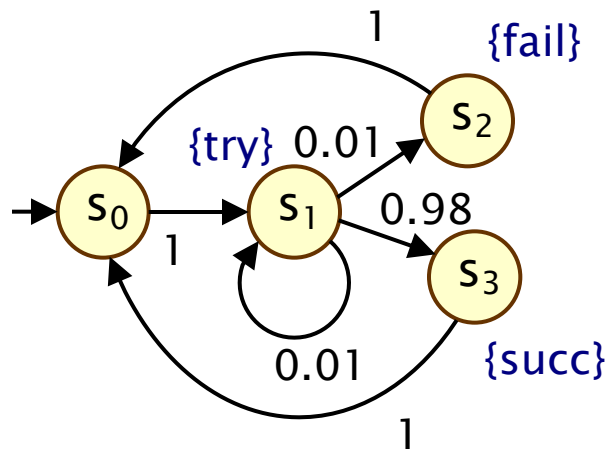
$$\mathbf{P} = \begin{bmatrix} 0 & 1 & 0 & 0 \\ 0 & 0.01 & 0.01 & 0.98 \\ 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \end{bmatrix}$$

$$\underline{x} \approx [ 0.332215, 0.335570, \\ 0.003356, 0.328859 ]$$

# Steady-state – Example

- Let  $\underline{x} = \underline{\pi}$
- Solve:  $\underline{x} \cdot \mathbf{P} = \underline{x}$ ,  $\sum_s \underline{x}(s) = 1$

$$\underline{x} \approx [ 0.332215, 0.335570, 0.003356, 0.328859 ]$$



$$\mathbf{P} = \begin{bmatrix} 0 & 1 & 0 & 0 \\ 0 & 0.01 & 0.01 & 0.98 \\ 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 \end{bmatrix}$$

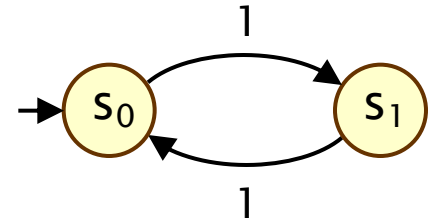
Long-run percentage of time spent in the state “try”  
 $\approx 33.6\%$

Long-run percentage of time spent in “fail”/”succ”  
 $\approx 0.003356 + 0.328859$   
 $\approx 33.2\%$

# Periodic DTMCs

- For (finite, irreducible) periodic DTMCs, this limit:

$$\underline{\pi}_s(s') = \lim_{k \rightarrow \infty} \underline{\pi}_{s,k}(s')$$



- in general does not exist, but this limit does:

$$\lim_{n \rightarrow \infty} \frac{1}{n} \cdot \sum_{k=1}^n \underline{\pi}_{s,k}(s')$$

(and where both limits exist,  
e.g. for aperiodic DTMCs,  
these 2 limits coincide)

- Steady-state probabilities for periodic DTMCs can still be computed, again by solving the same set of linear equations:

$$\underline{\pi} \cdot \mathbf{P} = \underline{\pi} \quad \text{and} \quad \sum_{s \in S} \underline{\pi}(s) = 1$$

# Steady-state – General case

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- General case: reducible DTMC
- there are multiple solutions of steady-state equation

$$\underline{\pi} \cdot \mathbf{P} = \underline{\pi} \quad \text{and} \quad \sum_{s \in S} \underline{\pi}(s) = 1$$

- number of (lin. Independent) solutions = number of BSCCs
- limiting distribution obtained by iterations exists
- limiting distribution depends on initial one

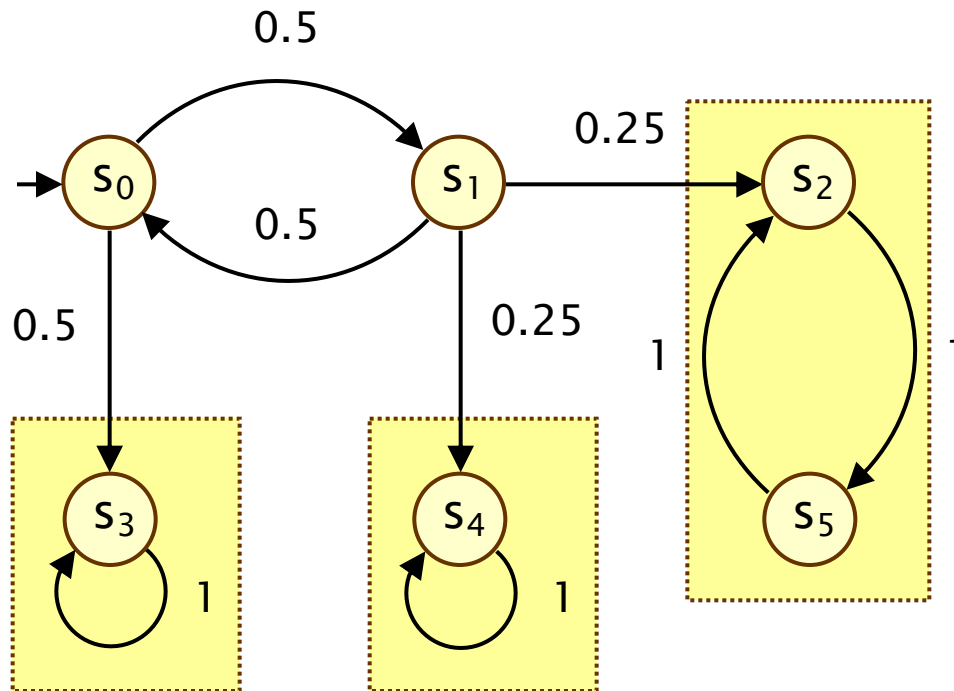
# Steady-state – General case

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- General case: reducible DTMC
  - compute vector  $\underline{\pi}_s$
- Compute BSCCs for DTMC; then two cases to consider:
- (1)  $s$  is in a BSCC  $T$ 
  - compute steady-state probabilities  $\underline{x}$  in sub-DTMC for  $T$
  - $\underline{\pi}_s(s') = \underline{x}(s')$  if  $s'$  in  $T$
  - $\underline{\pi}_s(s') = 0$  if  $s'$  not in  $T$
- (2)  $s$  is not in any BSCC
  - compute steady-state probabilities  $\underline{x}_T$  for sub-DTMC of each BSCC  $T$  and combine with reachability probabilities to BSCCs
  - $\underline{\pi}_s(s') = \text{ProbReach}(s, T) \cdot \underline{x}_T(s')$  if  $s'$  is in BSCC  $T$
  - $\underline{\pi}_s(s') = 0$  if  $s'$  is not in a BSCC

# Steady-state – Example 2

- $\underline{\pi}_s$  depends on initial state  $s$



$$\underline{\pi}_{s3} = [0 \ 0 \ 0 \ 1 \ 0 \ 0]$$

$$\underline{\pi}_{s4} = [0 \ 0 \ 0 \ 0 \ 1 \ 0]$$

$$\underline{\pi}_{s2} = \underline{\pi}_{s5} = \left[0, 0, \frac{1}{2}, 0, 0, \frac{1}{2}\right]$$

$$\underline{\pi}_{s0} = \left[0, 0, \frac{1}{12}, \frac{2}{3}, \frac{1}{6}, \frac{1}{12}\right]$$

$$\underline{\pi}_{s1} = \dots$$

# Qualitative properties

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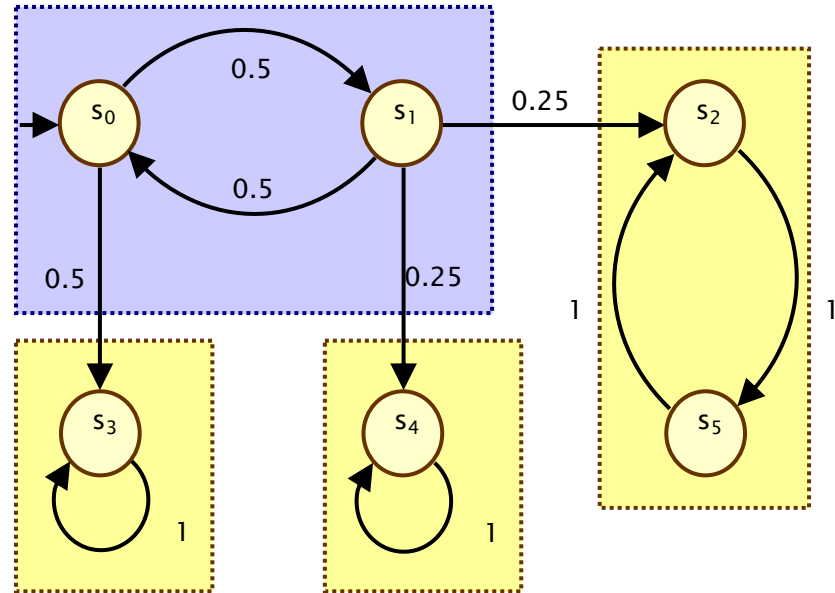
- **Quantitative** properties:
  - “what is the probability of event A?”
- **Qualitative** properties:
  - “the probability of event A is 1” (“**almost surely** A”)
  - or: “the probability of event A is  $> 0$ ” (“**possibly** A”)
- For finite DTMCs, qualitative properties do not depend on the transition probabilities – only need underlying graph
  - e.g. to determine “is target set T reached with probability 1?” (more in the DTMC model checking lecture later)
  - computing BSCCs of a DTMCs yields information about long-run qualitative properties...



# Fundamental property

- Fundamental property of finite DTMCs...

- With probability 1, a BSCC will be reached and all of its states visited infinitely often

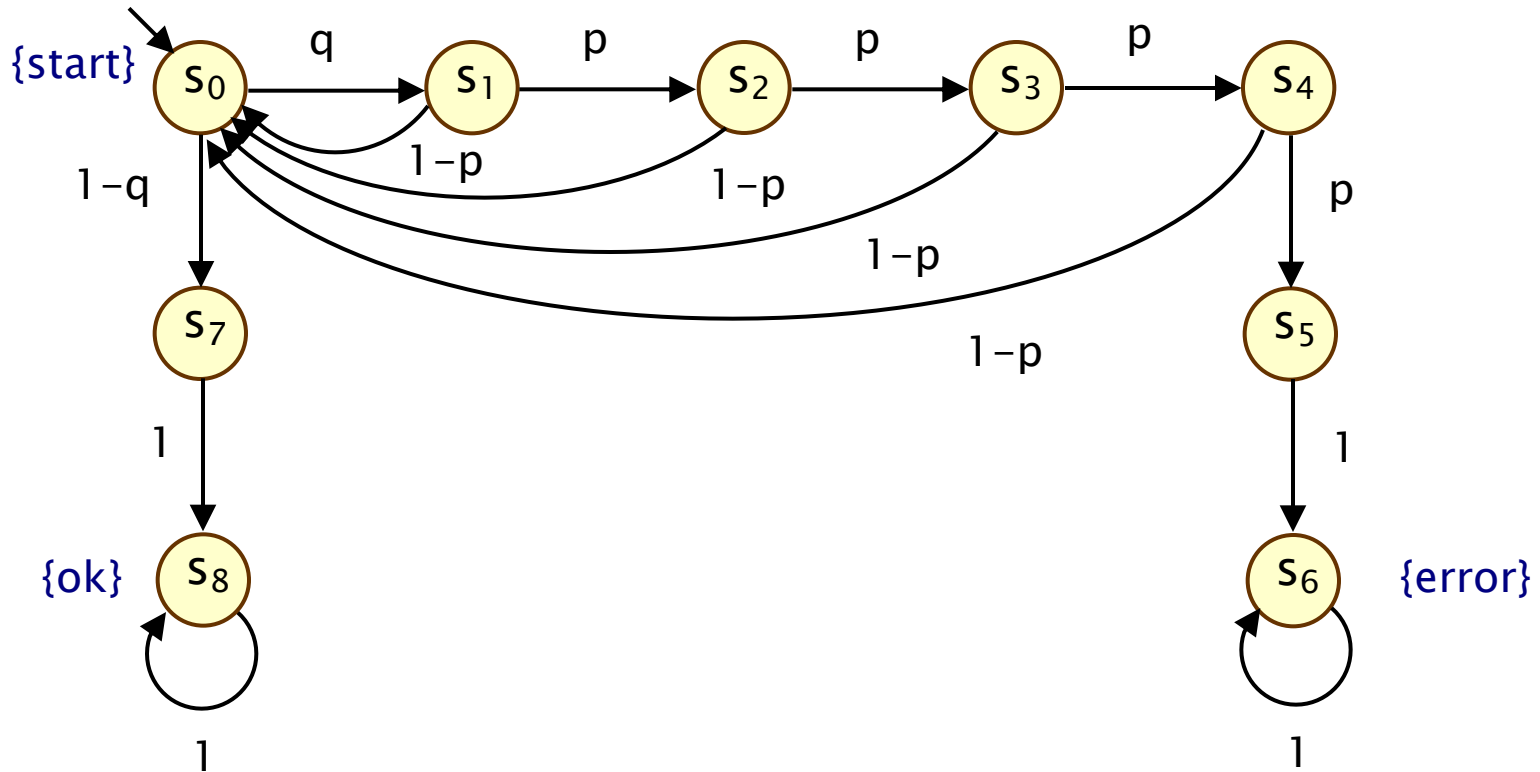


- Formally:

$$\begin{aligned}
 & - \Pr_{s_0} ( s_0 s_1 s_2 \dots \mid \exists i \geq 0, \exists \text{ BSCC } T \text{ such that} \\
 & \quad \forall j \geq i \ s_j \in T \text{ and} \\
 & \quad \forall s \in T \ s_k = s \text{ for infinitely many } k ) = 1
 \end{aligned}$$

# Zeroconf example

- 2 BSCCs:  $\{s_6\}$ ,  $\{s_8\}$
- Thus, probability of trying to acquire a new address infinitely often (i.e., visiting  $\{start\}$  i.o.) is 0



# Repeated reachability

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- Repeated reachability:  $GF B$ 
  - “always eventually...”, “infinitely often...”
- $\Pr_{s_0} ( s_0s_1s_2\dots \mid \forall i \geq 0 \exists j \geq i s_j \in B )$ 
  - where  $B \subseteq S$  is a set of states
- e.g. “what is the probability that the protocol successfully sends a message infinitely often?”
- Is this measurable? Yes...
  - set of satisfying paths is:  $\bigcap_{n \geq 0} \bigcup_{m \geq n} C_m$
  - where  $C_m$  is the union of all cylinder sets  $\text{Cyl}(s_0s_1\dots s_m)$  for finite paths  $s_0s_1\dots s_m$  such that  $s_m \in B$

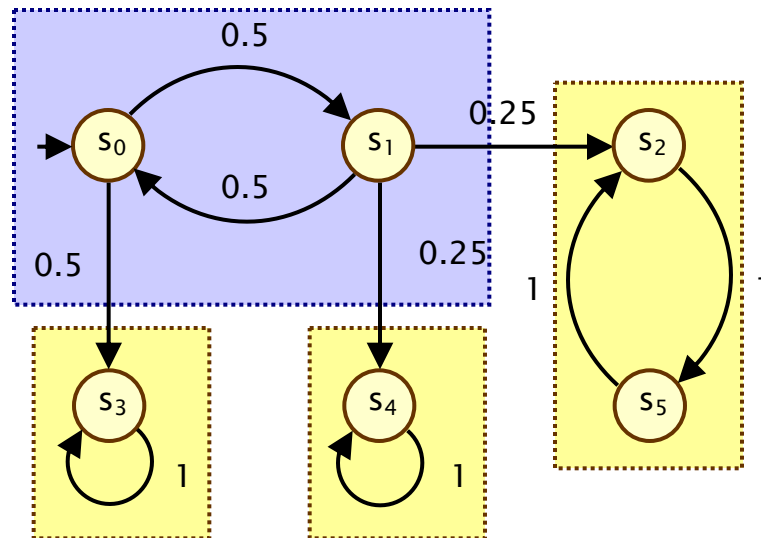
# Qualitative repeated reachability

- $\Pr_{s_0} ( s_0 s_1 s_2 \dots \mid \forall i \geq 0 \exists j \geq i s_j \in B ) = 1$   
 $\Pr_{s_0} ( \text{“always eventually B”} ) = 1$

if and only if

- $T \cap B \neq \emptyset$  for each BSCC  $T$  that is reachable from  $s_0$

Example:  
 $B = \{ s_3, s_4, s_5 \}$



# Persistence

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- Persistence properties:
  - “eventually forever...”
- $\Pr_{s_0} ( s_0 s_1 s_2 \dots \mid \exists i \geq 0 \forall j \geq i s_j \in B )$ 
  - where  $B \subseteq S$  is a set of states
- e.g. “what is the probability of the leader election algorithm reaching, and staying in, a stable state?”
- e.g. “what is the probability that an irrecoverable error occurs?”
- Is this measurable? Yes...

# Persistence

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- Persistence properties:
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- $\Pr_{s_0} ( s_0 s_1 s_2 \dots \mid \exists i \geq 0 \forall j \geq i s_j \in B )$ 
  - where  $B \subseteq S$  is a set of states
- e.g. “what is the probability of the leader election algorithm reaching, and staying in, a stable state?”
- e.g. “what is the probability that an irrecoverable error occurs?”
- Is this measurable? Yes...  $FG B = \neg GF (S \setminus B)$

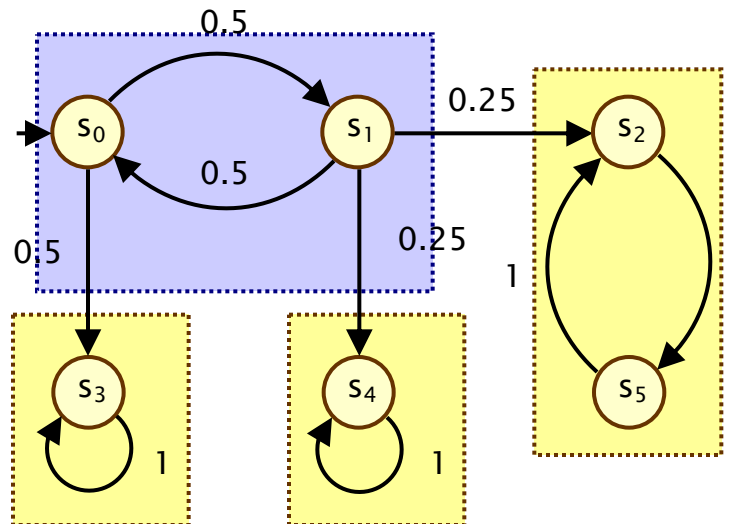
# Qualitative persistence

- $\Pr_{s_0} ( s_0 s_1 s_2 \dots \mid \exists i \geq 0 \forall j \geq i s_j \in B ) = 1$   
 $\Pr_{s_0} ( \text{“eventually forever B”} ) = 1$

if and only if

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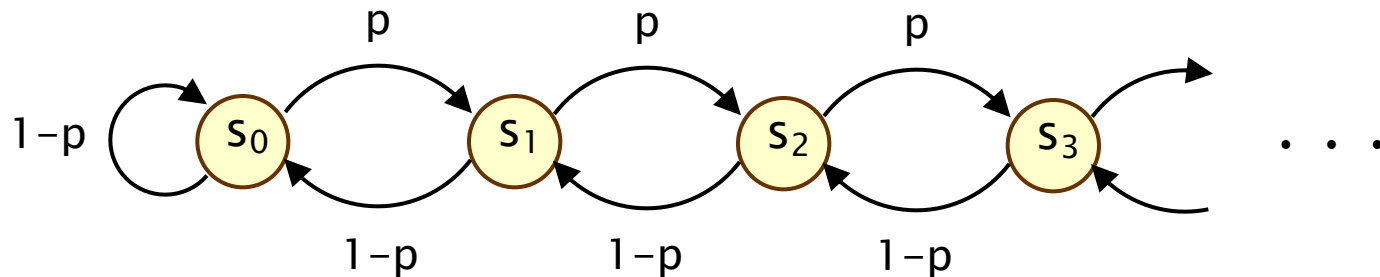
Example:  
 $B = \{ s_2, s_3, s_4, s_5 \}$



# Aside: Infinite-state Markov chains

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- Infinite-state random walk



- Value of probability  $p$  **does** affect qualitative properties
  - $\text{ProbReach}(s, \{s_0\}) = 1$  if  $p \leq 0.5$
  - $\text{ProbReach}(s, \{s_0\}) < 1$  if  $p > 0.5$
- (not comprehensively studied in this course)



# Summing up...

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- **Transient state probabilities**
  - successive vector–matrix multiplications
- **Long–run/steady–state probabilities**
  - requires graph analysis
  - irreducible case: solve linear equation system
  - reducible case: steady–state for sub–DTMCs + reachability
- **Qualitative properties**
  - repeated reachability
  - persistence