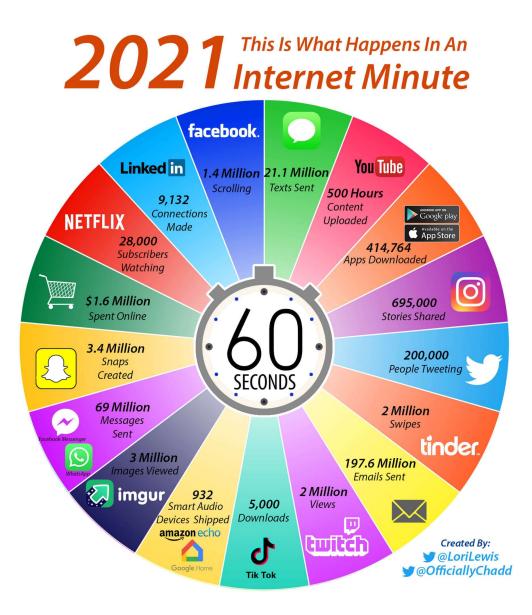
# Distributed platforms for Big Data

Prof. Carlo Ferrari Michele Stecca, Ph.D. a.y. 2023-2024

#### Data, data, data

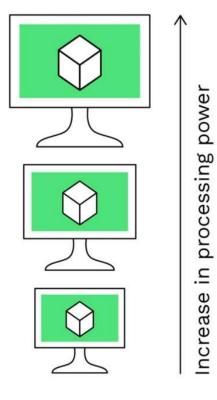


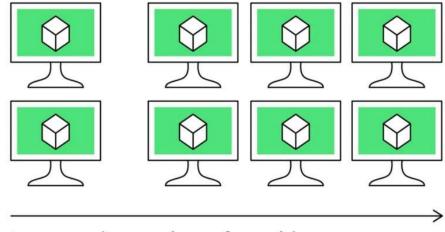
### Vertical vs. Horizontal scalability

#### Scalability

#### **Vertical scaling**

#### Horizontal scaling





Increase in number of machines

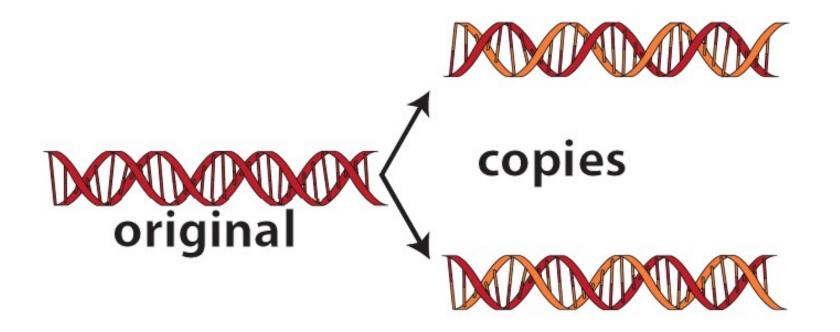
# Design principles

# (Data) Partitioning



## Design principles

# (Data) Replication

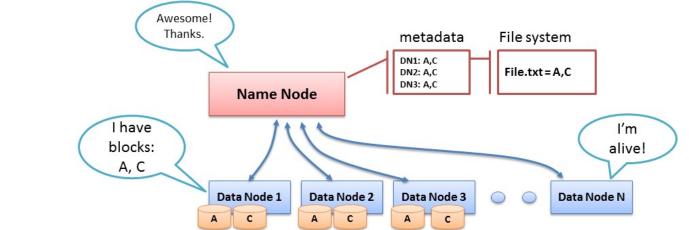


### An architecture for Big Data

#### **Distributed Filesystem**

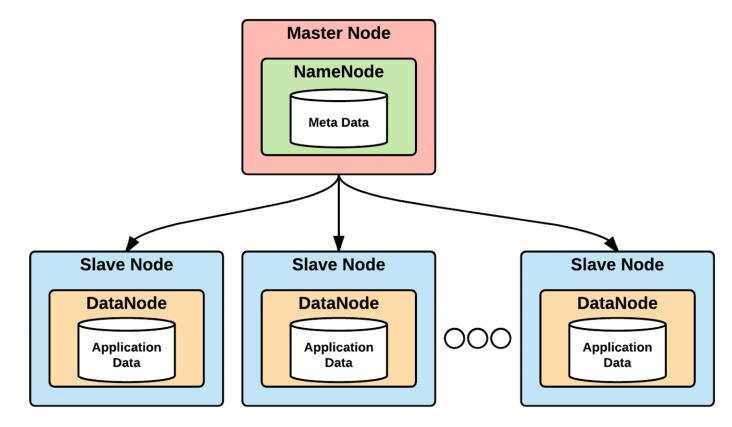


#### HaDoop FileSystem Architecture

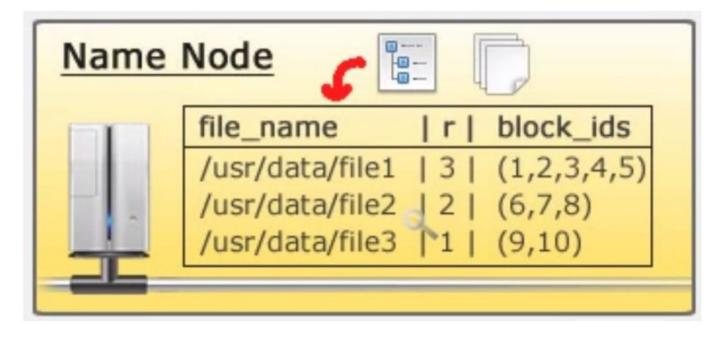


- Files are splitted into >64MB blocks
- Throughput optimized
- Write once read many
- Failure resistant: manages failures using block replication

Master & Slave architecture

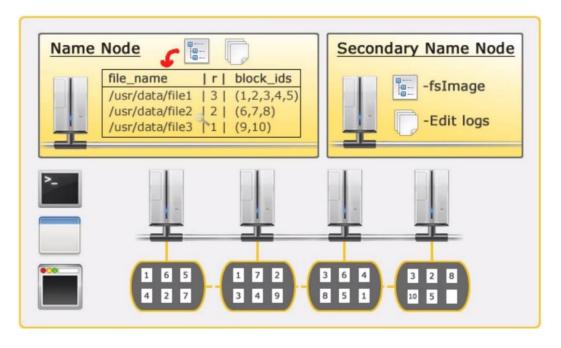


#### HaDoop FileSystem Architecture: Name Node



- Cluster's SPOF!
- Manages the file system
- Lists the files and blocks in which they are divided
- Manages strategies replication and block allocation
- Checks nodes's reliability

#### HaDoop FileSystem Architecture: DataNode



- Manages the storage and the client's requests
- Sends Heartbeat to NameNod

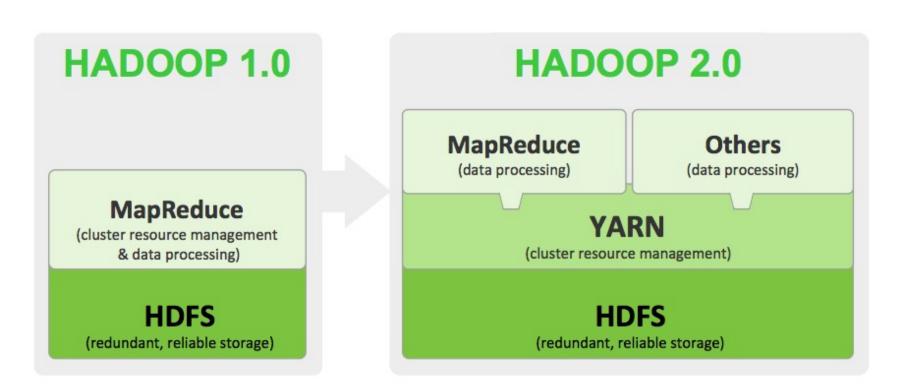
### An architecture for Big Data

#### Move computation close to data!!!

Distributed Computational Model + Execution Engine

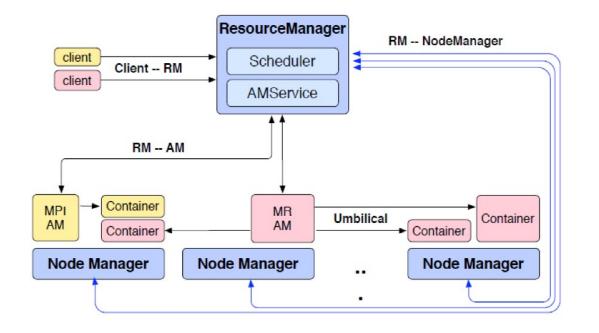


#### Hadoop versions



#### Cluster management: Yarn

Yet Another Resource Negotiator architecture:



#### Resource Manager:

- One per cluster global view
- No static resource partitioning
- Handle Job request
- Find a container to Application Manager

### An architecture for Big Data

Distributed Computational Model + Execution Engine: Map Reduce

Resource Manager (YARN)



### An architecture for Big Data

Applications («Pure» MR Apps, SQL, Machine Learning, Graphs, Streaming, etc.)

Distributed Computational Model + Execution Engine (Map Reduce)

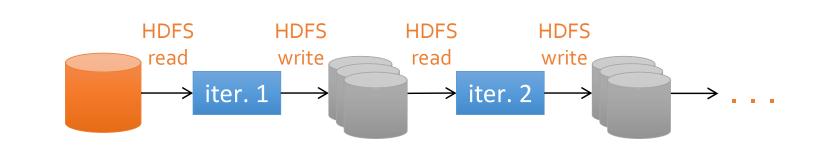
Resource Manager (YARN)



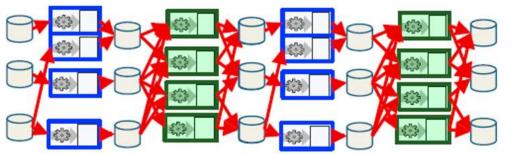
# Map Reduce limitations

- Not so flexible from a programmer point of view
- Not so efficient

. . .



• Iterative jobs involve a lot of disk I/O for each repetition





Applications («Pure» MR Apps, SQL, Machine Learning, Graphs, Streaming, etc.)

> Distributed Computational Model + Execution Engine: Map Reduce

Resource Manager (YARN)



Applications («Pure» Spark Apps, SQL, Machine Learning, Graphs, Streaming, etc.)

Distributed Computational Model + Execution Engine: Spark

Resource Manager (YARN)



Applications (Simple APIs, SQL, Machine Learning, Graphs, Streaming, etc.)

> Distributed Computational Model + Execution Engine: Spark

Resource Manager (YARN)

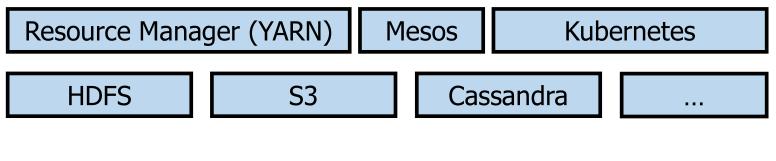
Mesos

Kubernetes



Applications (Simple APIs, SQL, Machine Learning, Graphs, Streaming, etc.)

> Distributed Computational Model + Execution Engine: Spark

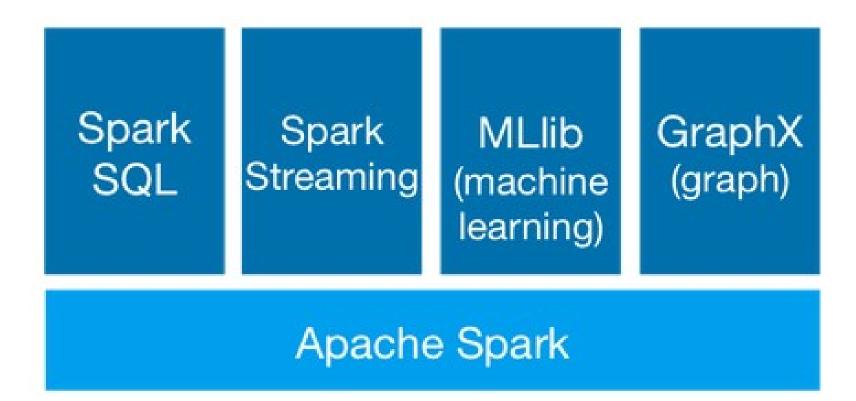




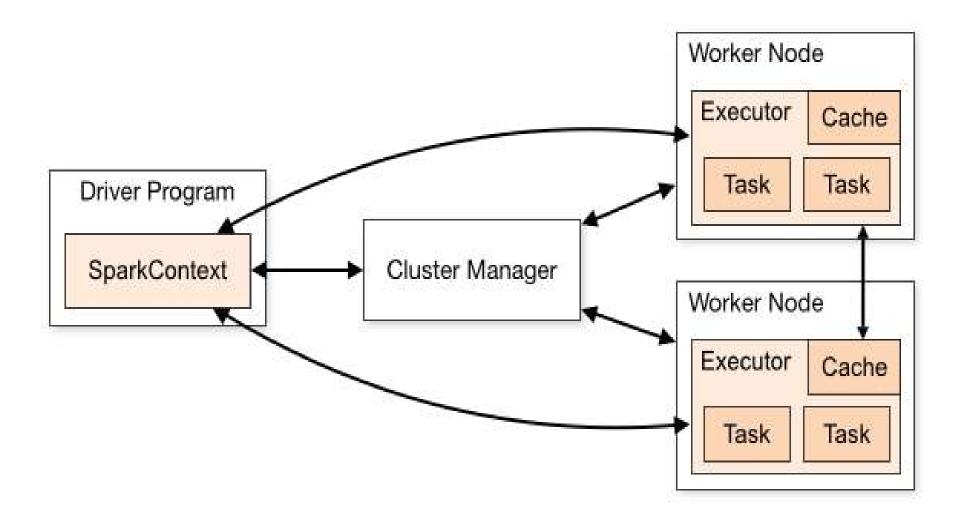
# MapReduce vs. Spark

- Spark is in-memory
- Less expensive shuffles
- There are many more primitives
- It supports Java, Scala, Python & R
- Interactive shells are available
- Generalized patterns
  - unified engine for many use cases
- Lazy evaluation of the lineage graph
  - reduces wait states, better pipelining
- Lower overhead for starting jobs

### Spark ecosystem



### Spark (Deployment)



### RDDs

- Resilient Distributed Datasets (RDDs) are the primary abstraction in Spark a fault-tolerant collection of elements that can be operated on in parallel
- 2 types of operations on RDDs:
  - transformations and actions
    - transformations are lazy (not computed immediately)
- however, an RDD can be *persisted* into storage in memory or disk

transformation	description
<pre>map(func)</pre>	return a new distributed dataset formed by passing each element of the source through a function <i>func</i>
<pre>filter(func)</pre>	return a new dataset formed by selecting those elements of the source on which <i>func</i> returns true
<pre>flatMap(func)</pre>	similar to map, but each input item can be mapped to 0 or more output items (so <i>func</i> should return a Seq rather than a single item)
<pre>sample(withReplacement, fraction, seed)</pre>	sample a fraction <i>fraction</i> of the data, with or without replacement, using a given random number generator seed
union(otherDataset)	return a new dataset that contains the union of the elements in the source dataset and the argument
<pre>distinct([numTasks]))</pre>	return a new dataset that contains the distinct elements of the source dataset

transformation	description
groupByKey([numTasks])	when called on a dataset of (K, V) pairs, returns a dataset of (K, Seq[V]) pairs
<pre>reduceByKey(func, [numTasks])</pre>	when called on a dataset of $(\kappa, v)$ pairs, returns a dataset of $(\kappa, v)$ pairs where the values for each key are aggregated using the given reduce function
<pre>sortByKey([ascending], [numTasks])</pre>	when called on a dataset of $(\kappa, v)$ pairs where $\kappa$ implements ordered, returns a dataset of $(\kappa, v)$ pairs sorted by keys in ascending or descending order, as specified in the boolean ascending argument
<pre>join(otherDataset, [numTasks])</pre>	when called on datasets of type $(\kappa, v)$ and $(\kappa, w)$ , returns a dataset of $(\kappa, (v, w))$ pairs with all pairs of elements for each key
<pre>cogroup(otherDataset, [numTasks])</pre>	when called on datasets of type (K, V) and (K, W), returns a dataset of (K, Seg[V], Seg[W]) tuples – also called groupWith
cartesian(otherDataset)	when called on datasets of types $\tau$ and $\upsilon$ , returns a dataset of $(\tau, \upsilon)$ pairs (all pairs of elements)

action	description
reduce(func)	aggregate the elements of the dataset using a function <i>func</i> (which takes two arguments and returns one), and should also be commutative and associative so that it can be computed correctly in parallel
collect()	return all the elements of the dataset as an array at the driver program – usually useful after a filter or other operation that returns a sufficiently small subset of the data
count()	return the number of elements in the dataset
first()	return the first element of the dataset – similar to take(1)
take(n)	return an array with the first <i>n</i> elements of the dataset – currently not executed in parallel, instead the driver program computes all the elements
<pre>takeSample(withReplacement, fraction, seed)</pre>	return an array with a random sample of <i>num</i> elements of the dataset, with or without replacement, using the given random number generator seed

action	description
<pre>saveAsTextFile(path)</pre>	write the elements of the dataset as a text file (or set of text files) in a given directory in the local filesystem, HDFS or any other Hadoop-supported file system. Spark will call tostring on each element to convert it to a line of text in the file
<pre>saveAsSequenceFile(path)</pre>	write the elements of the dataset as a Hadoop sequenceFile in a given path in the local filesystem, HDFS or any other Hadoop-supported file system. Only available on RDDs of key-value pairs that either implement Hadoop's writable interface or are implicitly convertible to writable (Spark includes conversions for basic types like Int, Double, String, etc).
countByKey()	only available on RDDs of type $(\kappa, v)$ . Returns a `Map` of $(\kappa, int)$ pairs with the count of each key
<pre>foreach(func)</pre>	run a function <i>func</i> on each element of the dataset – usually done for side effects such as updating an accumulator variable or interacting with external storage systems

- Spark can *persist* (or cache) a dataset in memory across operations
- Each node stores in memory any slices of it that it computes and reuses them in other actions on that dataset often making future actions more than 10x faster
- The cache is *fault-tolerant*: if any partition of an RDD is lost, it will automatically be recomputed using the transformations that originally created it



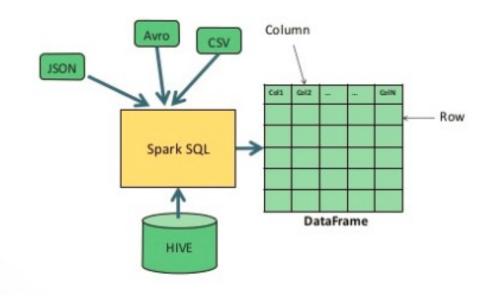
- Spark SQL is a Spark module for structured data processing
- Uses more information about the structure of both the data and the computation being performed
- Spark SQL uses this extra information to perform extra optimizations
- Integrated with Hive metastore

#### Spark SQL: the Dataframe abstraction

- The **DataFrame API** provides a **higher-level abstraction**, allowing you to use a query language to manipulate data. In fact, you can use **SQL**, as well.
- This code does essentially the same thing the previous RDD code does. Look how much easier it is to read.
- You have probably met DataFrames already in Python or R

#### Spark SQL: the Dataframe abstraction

- It provides the benefits of RDDs (strong typing, ability to use powerful lambda functions) with the benefits of Spark SQL's optimized execution engine
- It is conceptually equivalent to a table in a relational database or a data frame in R/Python, but with richer optimizations under the hood
- DataFrames can be constructed from a wide array of sources such as: structured data files, tables in Hive, external databases, or existing RDDs



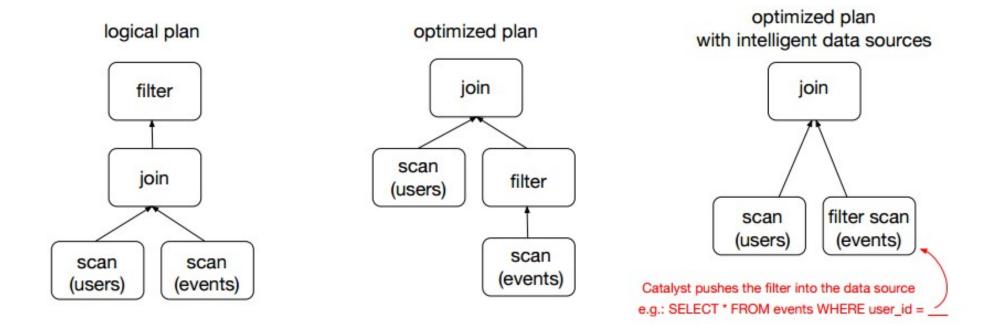
Spark SQL: the Dataframe abstraction

# Data Sources supported by DataFrames



# **Example Optimization**

users.join(events, users("id") === events("uid"))
 .filter(events("date") > "2015-01-01")



### Spark SQL: JOIN Operations

#### (Distributed) JOIN Types

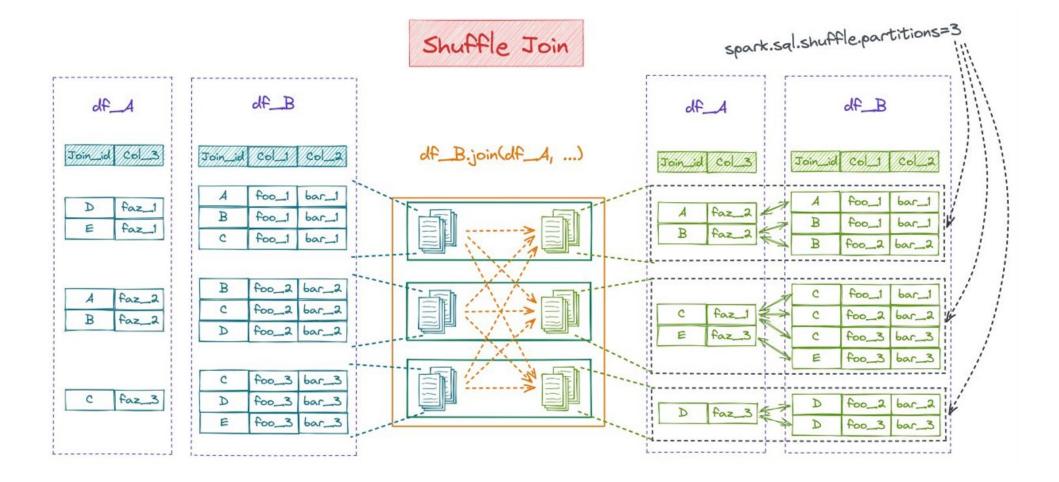
- Shuffle JOIN
- Broadcast JOIN
- Merge Sort JOIN
- Skew JOIN

#### **JOIN Optimization parameters**

**spark.sql.adaptive.enabled** - if this option is set to True Spark will make use of the runtime statistics to choose the most efficient query execution plan, one of the optimizations is automated conversion of shuffle join to a broadcast join.

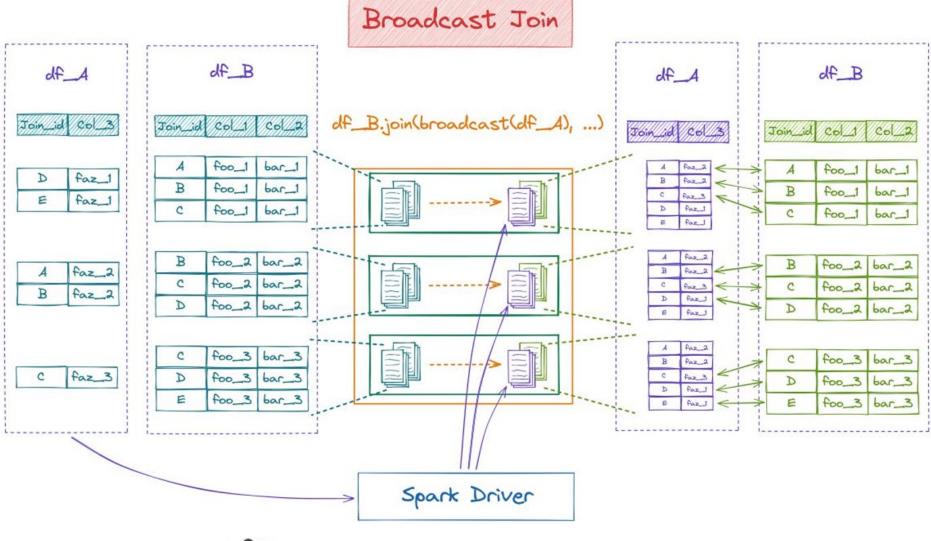
**spark.sql.autoBroadcastJoinThreshold** - denotes the maximum size of a dataset that would be automatically broadcasted.

### Spark SQL: JOIN Operations



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### Spark SQL: JOIN Operations

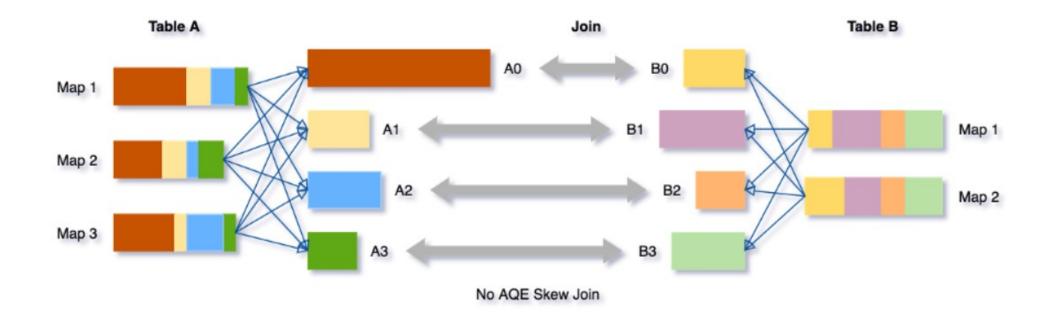




linkedin.com/in/aurimas-griciunas

#### Spark SQL: JOIN Operations (Data Skew case)

- Data skew is a condition in which a table's data is unevenly distributed among partitions in the cluster.
- Data skew can severely downgrade performance of queries, especially those with joins.
- Joins between big tables require shuffling data and the skew can lead to an extreme imbalance of work in the cluster.



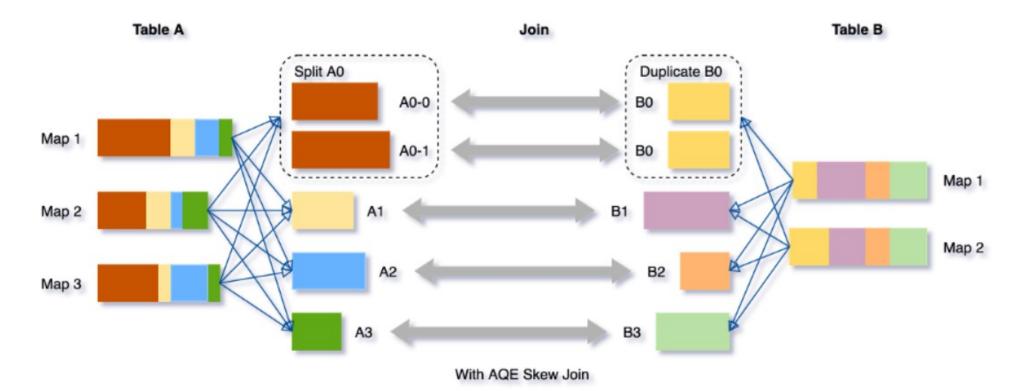
### Spark SQL: JOIN Operations (Data Skew case)

#### **JOIN Optimization parameters**

#### spark.sql.adaptive.optimizeSkewsInRebalancePartitions.enabled

- When true and spark.sql.adaptive.enabled is true, Spark will optimize the skewed shuffle partitions in RebalancePartitions and split them to smaller ones

AQE mechanisms transparently discover and optimize implementation.

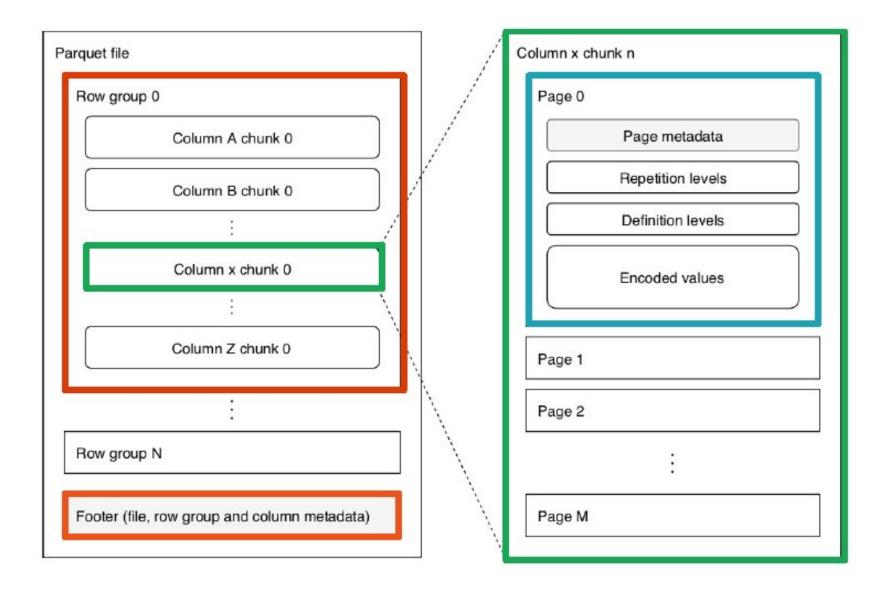


#### About data formats

# **Different workloads**

- OLTP
  - Online transaction processing
  - Lots of small operations involving whole rows
- OLAP
  - Online analytical processing
  - Few large operations involving *subset of all columns*

#### The Parquet (hybrid) file format



# **Optimization: predicate pushdown**

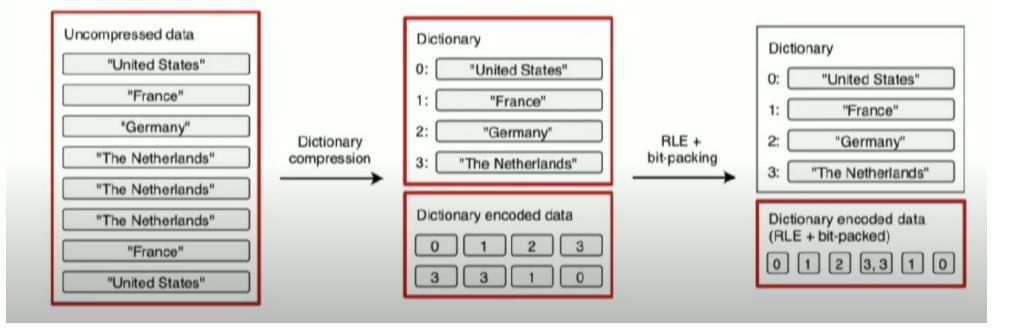
SELECT \* FROM table WHERE x > 5

Row-group 0: x: [min: 0, max: 9] Row-group 1: x: [min: 3, max: 7] Row-group 2: x: [min: 1, max: 4]

...

# **Parquet: encoding schemes**

• RLE\_DICTIONARY



#### References

#### https://spark.apache.org/

Performance Tuning: https://spark.apache.org/docs/latest/sql-performance-tuning.html

https://www.agilelab.it/blog/spark-3-0-first-hands-on-approach-with-adaptive-query-execution-part-3