Master Degree in Computer Engineering

Final Exam for Automata, Languages and Computation

September 3rd, 2021

- 1. [4 points] Consider the regular expression $r = (0+1)^* \emptyset(\epsilon+01)$. Convert r into an equivalent ϵ -NFA using the construction we have presented in class. Important: do not simplify the regular expression before applying the construction, use r as is.
- 2. [9 points] Consider the following languages, defined over the alphabet $\Sigma = \{a, b, c\}$

$$L_{1} = \{ w \mid w = XuYvZ, X, Y, Z \in \Sigma, u, v \in \Sigma^{*}, X = Z, |u| = |v| \};$$

$$L_{2} = \{ w \mid w = XuYvZ, X, Y, Z \in \Sigma, u, v \in \Sigma^{*}, X = Y = Z \};$$

$$L_{3} = \{ w \mid w = XuYvZ, X, Y, Z \in \Sigma, u, v \in \Sigma^{*}, X = Y = Z, |u| = |v| \}.$$

State whether the above are regular languages, and provide a mathematical proof of your answers.

- 3. [6 points] Assess whether the following statements are true or false, providing motivations for all of your answers.
 - (a) If $L_1, L_2 \in CFL$ then $L_1 \cap L_2 \in REC$;
 - (b) If $L_1 \in \text{REG}$ and $L_2 \in \text{CFL} \setminus \text{REG}$ then $L_1 \cdot L_2$ is never in REG;
 - (c) If $L_1 \cdot L_2 \in \text{REG}$ then $L_1, L_2 \in \text{REG}$.

(please see next page)

4. [6 points] Consider the language $L = \{a^n b^m \mid n, m \ge 0\}$ and the context-free grammar $G = (\{S, A\}, \{a, b\}, P, S)$, where P contains the following rules

Using mutual induction, we want to construct a mathematical proof that L(G) = L. In order to do this, for each variable $X \in \{S, A\}$ define a property \mathcal{P}_X over strings $x \in \{a, b\}^*$ such that $\mathcal{P}_X(x)$ holds if and only if there is a derivation in G that starts with X and produces x. **Important**: develop your proof only for property \mathcal{P}_S .

5. [8 points] Define the notion of property of the languages generated by TMs and state Rice's theorem. Provide the proof of Rice's theorem that we have developed in class.