

LECTURE 20 , May 18, 2023

GENERAL THEORY of PDEs for DIFFERENTIAL GAMES?

- NOT for Non-zero sum games!
because systems of 1st. order nonlinear
are NASTY.
- YES for 0-sum games & the Isaacs eqn

PLAN

- def. a value function.
- DYNAMIC PROGRAMMING PRINCIPLE
- value is a visc. sol. of Isaacs eq. ...
.... THE UNIQUE ONE.

Q How to define the value function?

N.B. open loop vs. over loop is not good!

History: • R. Isaacs 50-60 ! method of characteristics

... piecewise C^1 solutions. $n \leq 2$

• W. Fleming (1961)

$$\overbrace{\quad\quad\quad\quad}^{\Delta t} +$$

Euler scheme, solve static

series at discrete times, let $\Delta t \rightarrow 0$.

- A. Friedman , time discrete $\Delta t \rightarrow 0$...
- N. Krasovskii - A. Subbotin 1969 - 1977

"Positional differential games."

also involves a limit procedure.

• NON-ANTICIPATING (CAUSAL or MARKOVIAN) STRATEGIES

Varaiya, Roxin, Elliott-Kalton (1967 - 72)

Def: L.C Evans - Souganidis 1984 \leftrightarrow viscosity sols.

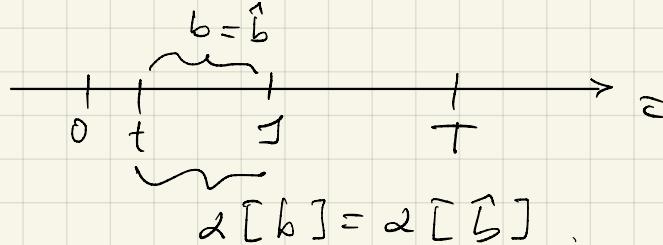
Notations: $\Omega_t := \{ \alpha : [t, T] \rightarrow A \text{ meas. le} \}$.

$\beta_t := \{ b : [t, T] \rightarrow B \text{ meas. le} \}$.

Def. A strategy of 1st player is $\alpha : \beta_t \rightarrow \Omega_t$, i's

NONANTICIPATING if. $\forall t \leq s \leq T, \forall b, \tilde{b} \in \beta_t, b(\tau) = \tilde{b}(\tau) \quad \forall t \leq \tau \leq s$

$$b(\tau) = \tilde{b}(\tau) \quad \forall t \leq \tau \leq s \Rightarrow \alpha[b](\tau) = \alpha[\tilde{b}](\tau) \quad \forall t \leq \tau \leq s$$



$T_t = \{ \text{nonantic. strat. 1}^{\text{st}} \text{ player.} \}$,

$$\Delta_t = \{ \alpha \in 2^{\omega} \} \ni \beta$$

$\beta : \Omega_t \rightarrow \beta_t, \forall s \in [t, T], \forall \alpha, \hat{\alpha} \in \Omega_t :$

$$\alpha(\tau) = \hat{\alpha}(\tau) \quad \forall \tau \in [t, s] \Rightarrow \beta[\alpha](\tau) = \beta[\hat{\alpha}](\tau) \quad \forall \tau \in [t, s]$$

N.B. $\forall b \in \beta_t \quad \forall \alpha \in T_t \quad \exists \text{ unique trajectory of}$

$$\begin{cases} \dot{y}(s) = f(y(s), \alpha[b](s), b(s)) & s \geq t \\ y(t) = x \end{cases}$$

define $y(s) = y_x(s; t, \alpha[b], b)$.

Similarly $\alpha \in \Omega_t$, $\beta \in \Delta_t$ $\dot{y} = f(y, \alpha, \beta[\epsilon])$.

Payoff (for 1st) - cost (for 2nd) functional

$$J(x, t; \alpha(\cdot), b(\cdot)) = \int_t^T \ell(y_x(s), \alpha(s), b(s)) ds + g(y_x(T))$$

DEF. The LOWER VALUE of the D.C. is

$$V(t, x) := \inf_{\beta \in \Delta_t} \sup_{\alpha \in \Omega_t} J(x, t, \alpha, \beta[\epsilon]),$$

$$U(t, x) := \sup_{\alpha \in \Gamma_t} \inf_{b \in \beta_t} J(x, t, \alpha[b], b).$$

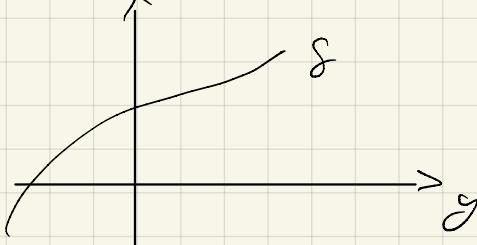
If $V(t, x) = U(t, x)$ the D.C. has a value.

Q: $V \leq U$?

Example (Berkovits) $\begin{cases} \dot{y} = (a - b)^2 \geq 0 \\ y(t) = x \end{cases}$ in \mathbb{R}

$$A = B = \{0, 1\}, \ell \equiv 0, g' > 0$$

b minimizes: she with $\dot{y} = 0$



$$\beta^*[\epsilon](s) = a(s) \Leftrightarrow \beta^* \in \Delta_t, \text{ s.t. } \dot{y}^* = 0 \text{ A.s.}$$

$$\forall \alpha \in \Omega_t, y^*(s) = x \Leftrightarrow \alpha \in \Omega_t$$

$$\Rightarrow V(t, x) \leq \sup_{\alpha \in \Omega_t} g(y^*(T)) = g(x)$$

• MAXIMISSES, she wants $\dot{y} = 1$

$$\alpha^*[b](s) := \begin{cases} 1 & \text{if } b(s) = 0 \\ 0 & \text{if } b(s) = 1 \end{cases} \quad \alpha^* \in \mathcal{P}_t$$

$$\dot{y}^*(s) = (\alpha^*[b](s) - b(s))^2 = 1 \quad \forall s \quad \forall b \in \mathcal{B}_t$$

$$\Rightarrow y^*(s) = x + s - t$$

$$U(t, x) \geq \inf_{b \in \mathcal{B}_t} g(y^*(t)) = g(x + t - t) \xrightarrow[t \geq 0]{} g(x) > g(x)$$

$$\Rightarrow V(t, x) < U(t, x) \quad \forall x \quad \forall t < T. \quad \blacksquare$$

Example of STRATEGIES.

Ex. 1. CONSTANT STRAT. Fix $\bar{b} \in \mathcal{B}_t$

$\beta[a] = \bar{b} \quad \forall a \in \Delta_t \Rightarrow \beta \in \Delta_t \Rightarrow \Delta_t \supseteq \text{a copy of } \mathcal{B}_t$

Ex. 2. $\Psi : A \rightarrow B$, $\stackrel{\text{ass.}}{\exists} j \mapsto \Psi \circ \alpha(j)$ meas. le. $\forall a \in \Delta_t$

$\Rightarrow \beta[a](s) = \Psi(\alpha(s))$ is holomorphic.

Ex. 3. Feedback: $\Phi : \mathbb{R}^n \times [0, T] \rightarrow B$ s.t.

$$\forall a \in \Delta_t \quad \exists \text{ unique sol. of} \quad \left\{ \begin{array}{l} \dot{q}(s) = f(q(s), \alpha(s), \Phi(q(s), s)) \\ q(0) = x \end{array} \right. \quad s \geq t.$$

Is spse. $s \mapsto \Phi(q(s), s)$ meas. le. Then

$\beta[a](s) := \Phi(q(s), s)$ is holomorphic? si, $\beta \in \Delta_t$

$$\Phi y_x(s; t, \alpha, \beta[a]) = q(s). \quad \blacksquare$$

Rmk. Information pattern is not very realistic,
but all more realistic values $\in [\bar{V}, \bar{U}]$, so
if $\bar{V} = \bar{U}$ they all coincide.

Es. If $\bar{V} = \bar{U}$ & if saddle \bar{V} and admissible
feedbacks $\Rightarrow \bar{V} = \bar{U} = \bar{V}$. □

Assumptions. $T > 0$, A, B metric compact.

- $f: \mathbb{R}^n \times A \times B \rightarrow \mathbb{R}^n$ cont., $|f| \leq d_1$
 $|f(x, a, s) - f(\tilde{x}, a, s)| \leq d_1 |x - \tilde{x}| \quad \forall x, \tilde{x} \in \mathbb{R}^n \text{ a.s.t. } a \in A, b \in B$
- $\ell: \mathbb{R}^n \times A \times B \rightarrow \mathbb{R}$ cont., $|\ell| \leq d_2$
 $|\ell(x, a, s) - \ell(\tilde{x}, a, s)| \leq d_2 |x - \tilde{x}| \quad \forall x, \tilde{x} \in \mathbb{R}^n \text{ a.s.t. } a \in A, b \in B$
- $g: \mathbb{R}^n \rightarrow \mathbb{R}$, $|g| \leq d_3$, $|g(x) - g(\tilde{x})| \leq d_3 |x - \tilde{x}|$.

Rmk. I can add $\dot{y}_{n+1} = \ell(g, a, b)$

$\tilde{g}(y, y_{n+1}) = g(y) + y_{n+1}$ # set an equivalent
game with $\tilde{\ell} \equiv 0 \in \mathcal{E}_{\tilde{g}}^{\uparrow}$. □

DYNAMIC PROGRAMMING PRINCIPLE (TENET OF TRANSITION).

Thm. $0 \leq t < t+\sigma \leq T \quad \forall x \in \mathbb{R}^n$

$$V(t, x) = \inf_{B \subseteq A_T} \sup_{a \in Q_t} \left\{ \int_t^{t+\sigma} \ell(g(s), a(s), \rho[a](s)) ds + V(t+\sigma, g(t+\sigma)) \right\}$$

where $y(\sigma) = y_x(\sigma; t, \alpha, \beta[\sigma])$

$$V(t, x) = \sup_{\alpha \in \Omega_t} \inf_{b \in \mathcal{B}_t} \left\{ \int_t^{t+\alpha} \rho(y(\sigma), d[b](\sigma), b(\sigma)) + \theta(t+\sigma, y(t+\sigma)) \right\}$$

$$y(\sigma) = y_x(\sigma; t, \alpha, \beta[\sigma]).$$

Pf. only for $\ell = 0$, $\notin \mathcal{V}$, only " \leq " (" \geq " in the last optional)

$$\frac{1}{2} \text{ thesis: } V(t, x) \leq W(t, x) := \inf_{\beta} \sup_{\alpha} V(t+\alpha, y(t+\alpha))$$

$\forall \varepsilon > 0 \quad \exists \delta \in \Delta_t : \text{(def. of inf.)}$

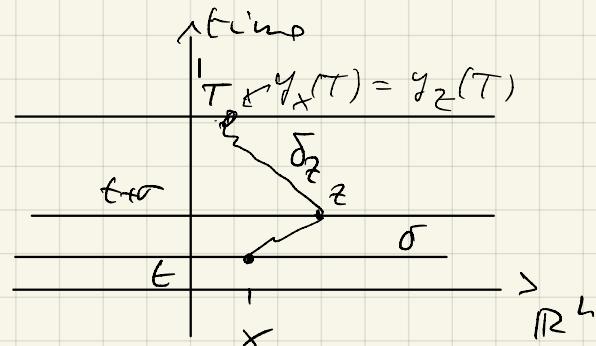
$$(1) \quad W(t, x) \geq \sup_{\alpha \in \Omega_t} V(t+\alpha, y_x(t+\alpha; t, \alpha, \delta[\alpha])) - \varepsilon$$

$$t \in \mathbb{R}^L \quad \text{def of inf.} \Rightarrow \exists \delta_2 \in \Delta_{t+\alpha}$$

$$(2) \quad V(t+\alpha, z) \geq \sup_{\alpha \in \Omega_{t+\alpha}} g(y_z(T; t+\alpha, \alpha, \delta_2[\alpha])) - \varepsilon$$

$$\text{Def. } \bar{\beta}[\alpha](\sigma) = \begin{cases} \delta[\alpha](\sigma) & t \leq \sigma \leq t+\alpha \\ \delta_2[\alpha](\sigma) & t+\alpha \leq \sigma \leq T \end{cases}$$

$$\text{with } z = y_x(t+\alpha; t, \alpha, \delta[\alpha])$$



$\bar{\beta} \in \Delta_t$? Yes, (1) & (2) \Rightarrow

$$W(t, x) \geq \sup_{\alpha \in \Omega_t} \sup_{\substack{z \in \Omega_{t+\alpha} \\ t \leq \sigma \leq t+\alpha}} g(y_z(T; t+\alpha, \alpha, \delta_2[\alpha])) - 2\varepsilon$$

$$\sup_{\alpha \in \Omega_t} g(y_x(T; t, \alpha, \bar{\beta}[\alpha]))$$

$$\geq \sup_{\alpha \in \Omega_t} g(\overbrace{\quad}^T) - 2\varepsilon \Rightarrow$$

$$W(t, x) \geq \inf_{\beta \in A_t} \sup_{a \in \partial_t} g(y_x(t; t, a, \beta[a])) \approx$$

$\overbrace{\quad \quad \quad}^{\mathcal{V}(t, x)}.$

Let $\varepsilon \rightarrow 0$ & set $W \geq \mathcal{V}$. $\forall t, x$ \blacksquare

Estimates $\Rightarrow \overbrace{\mathcal{V} \notin \mathcal{V}}$

Thm. $\exists C_4$ dep. on T, c_1, c_2, c_3 s.t.

$$1. \quad |\mathcal{V}(t, x)|, |\mathcal{V}(t, \hat{x})| \leq C_4 \quad \forall (t, x) \in [0, T] \times \mathbb{R}^L$$

$$2. \quad |\mathcal{V}(t, x) - \mathcal{V}(t, \hat{x})| \leq C_4 (|t - \hat{t}| + |x - \hat{x}|)$$

$$|\mathcal{V}(t, x) - \mathcal{V}(\hat{t}, \hat{x})| \leq \text{same}.$$

Part of proof : 2. $\left| \int_t^T l(c(s)) ds + g(c) \right| \leq (T-t) C_2 + C_3 \leq$

$$\leq T C_2 + C_3 =: C_4 \quad \blacksquare$$

2. optional, see Notes. \blacksquare

Isaacs' Hamiltonians

$$H^+(x, p) := \min_{b \in B} \max_{a \in A} \{ p \cdot f(x, a, b) + l(x, a, b) \}$$

$$H^-(x, p) := \max_{a \in A} \min_{b \in B} \{ \text{same} \} \leq H^+(x, p).$$

Prop. $H^+, H^- : \mathbb{R}^n \times \mathbb{R}^n \rightarrow \mathbb{R}$ are continuous and for $F = H^+ + H^-$

$\exists K \geq 0 : (L_x) |F(x, p) - F(\tilde{x}, p)| \leq K|x - \tilde{x}|(1 + |p|)$

$(L_p) |F(x, p) - F(x, \tilde{p})| \leq C_1 |p - \tilde{p}|$

Pf. 1. (L_x) for $F = H^+$ (H^- HW).

$H^+(\tilde{x}, p) - H^+(x, p) \leq \dots$ choose $b' \in B$:

$$H^+(x, p) = \max_a \{ p \cdot f(x, a, b') + l(x, a, b') \}$$

$$\geq p \cdot f(x, a, b') + l(x, a, b') \quad \forall a \in A.$$

$$H^+(\tilde{x}, p) \leq \max_a \{ p \cdot f(\tilde{x}, a, b') + l(\tilde{x}, a, b') \} \stackrel{\exists a'}{=} \underline{p \cdot f(\tilde{x}, a', b') + l(\tilde{x}, a', b')}$$

$$H^+(\tilde{x}, p) - H^+(x, p) \leq \downarrow + p \cdot f(x, a', b') + l(x, a', b')$$

$$\leq |p|C_1|x - \tilde{x}| + C_2|x - \tilde{x}| \leq K|x - \tilde{x}|(1 + |p|)$$

if $K := C_1 + C_2$. F is Lip $x \neq \tilde{x} \Rightarrow$

(L_x) for H^+ . \square

$\geq (L_p) HW$

3. $(L_x) + (L_p) \Rightarrow$ cont. of $H^+ + H^-$. \square