

Solutions to some exercises

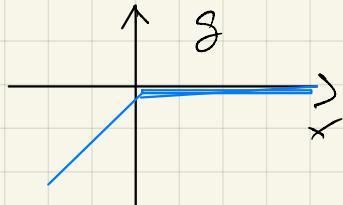
1. $u_t + \frac{(u_x)^2}{2} = 0$ in $\mathbb{R}^4 \times (0, \infty)$

a) compute u by H-L if $u(x, 0) = g(x) = -x^-$

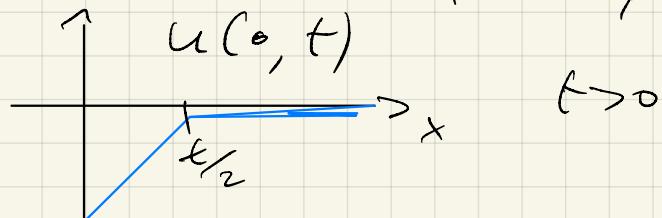
b) same for $g(x) = x^+$

c) $v = u_x$ $v_t + \left(\frac{v^2}{2}\right)_x = 0$

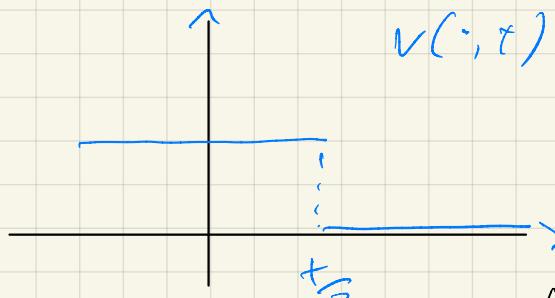
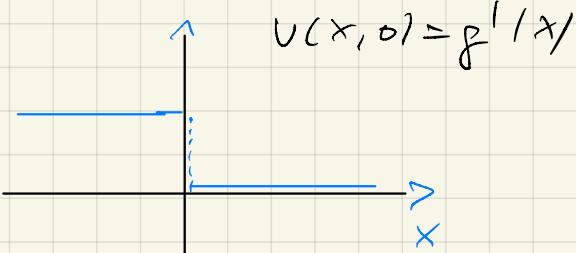
Draw the graphs!

Sol. a)

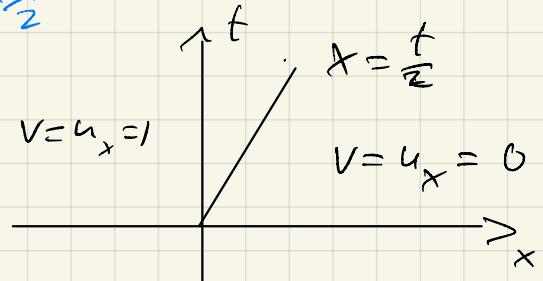
$$\text{H-L } u(x, t) = \begin{cases} x - \frac{t}{2}, & x < \frac{t}{2} \\ 0, & x > \frac{t}{2} \end{cases}$$



Comment & c)



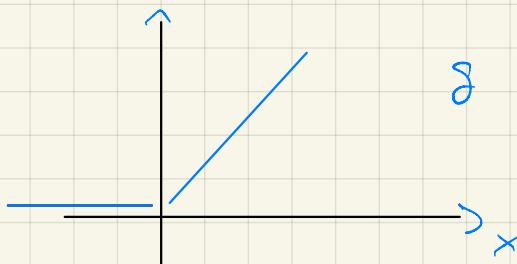
$$v(x, t) = u_x(x, t) = \begin{cases} 1, & x < \frac{t}{2} \\ 0, & x > \frac{t}{2} \end{cases}$$



b)

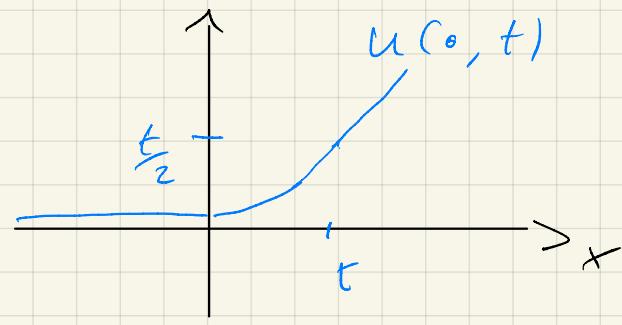
$$u(x, 0) = x^+$$

$\frac{1}{2} g(x)$



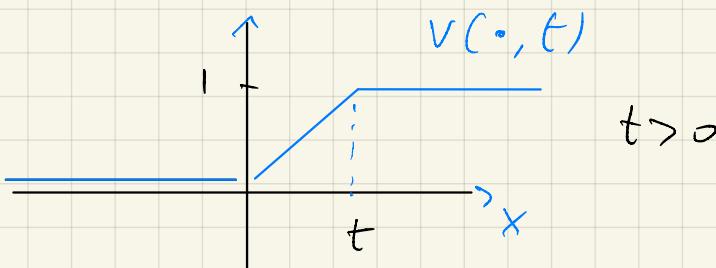
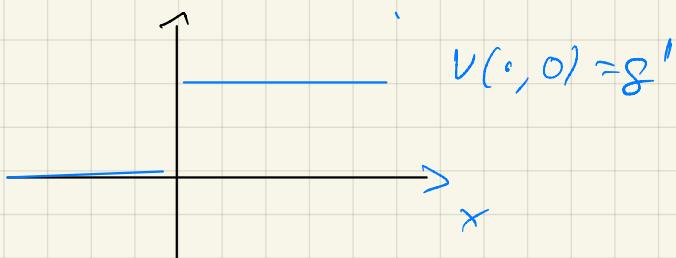
$$H-L \Rightarrow$$

$$u(x, t) = \begin{cases} 0 & x < 0 \\ x^2/2t, & 0 < x < t \\ x - \frac{t}{2} & x > t \end{cases}$$



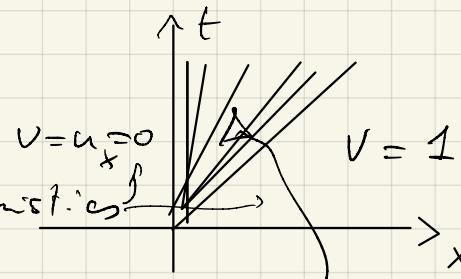
b-c) look at $v = u_x$

$$v(x, t) = \begin{cases} 0 & x < 0 \\ x/t & 0 < x < t \\ 1 & x > t \end{cases}$$



RAREFACTION WAVE

N.B.: from characteristics
we had no info. on v in



MIXED STRATEGIES for 0-SUM GAMES

A, B compact $\Phi_{\text{ec}}(A \times B)$

$i \in A$

$\mu \in \mathcal{P}(A) \iff \sum_{i \in A} \mu(i) \text{ Random variable with dist. } \mu$
 $\nu \in \mathcal{P}(B) \iff \sum_{j \in B} \nu(j) \text{ in } B$
 i.e. $P(\underline{X} \in S) = \mu(S)$, $P(\underline{Y} \in T) = \nu(T)$

S Borel set $\subseteq A$ $T \subseteq B$

Rmk.: $\tilde{\Phi}(\mu, \nu) = \iint_{A \times B} \Phi(a, b) d\mu(a) d\nu(b) =$

if $\underline{X}, \underline{Y}$ are independent. $= E[\Phi(\underline{X}, \underline{Y})] = \text{the expectation of } \Phi(\underline{X}, \underline{Y})$

$$\tilde{\Phi} : P(A) \times P(B) \rightarrow \mathbb{R}.$$

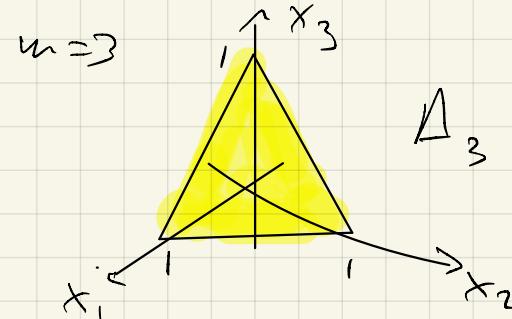
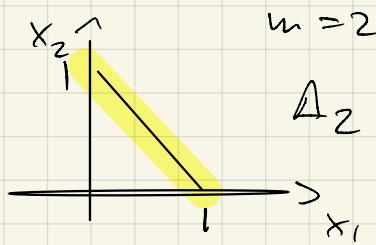
IMPORTANT EXAMPLE $A = \{1, \dots, m\}$, $B = \{1, \dots, n\}$.

$$\mu \in P(A) \quad \mu \leftrightarrow x = (x_1, \dots, x_m) \quad 0 \leq x_i \leq 1$$

$$x = \text{discrete density of } \mu \quad \mu(\{i\}) = x_i \quad \sum_{i=1}^m x_i = 1$$

$$y \leftrightarrow y = (y_1, \dots, y_n)$$

$$P(A) \iff \Delta_m = m\text{-dimensional simplex} = \\ = \{x \in [0, 1]^m : \sum_{i=1}^m x_i = 1\}.$$



Matrix game $\leftrightarrow M = (\phi_{ij})$ $\phi_{ij} = \tilde{\Phi}(i, j)$

$$\tilde{\Phi} : \Delta_m \times \Delta_n \rightarrow \mathbb{R}. \quad \tilde{\Phi}(x, y) = \sum_{\substack{i=1, \dots, m \\ j=1, \dots, n}} \phi_{ij} x_i y_j = x^T M y$$

Cor (Von Neumann) In 2-person 0-sum MATRIX GAMES,

the value in mixed strategies exists.

(i.e. the game $(P(A), P(B), \tilde{\Phi})$ has a value)

Pf. Δ_m, Δ_n are compact convex, $\tilde{\Phi}(x, y) = x^T M y$ is cont. & CONCAVE-CONVEX. V.N.Thm $\Rightarrow \tilde{v}^+ = \tilde{v}^-$

& \exists saddle point $(x^*, y^*) \leftrightarrow (\mu^*, \nu^*)$. ■

COMPUTATION OF SADDLE POINTS IN MIXED STRATEGIES

Ex HW: If (A, B, Φ) has a value v in pure strat.

$$\Rightarrow \tilde{v}^+ = v = \tilde{v}^- \Rightarrow \tilde{v} = v \quad \text{by}$$

Look for $(x^*, y^*) \in \Delta_m \times \Delta_n$ saddle for $x^T M y = \tilde{\Phi}$,

i.e. x^* is a SECUR. STRAT. for Δ_m , y^* is S.S. for Δ_n , i.e.

$$\tilde{\Phi}^{\min}(x^*) = \max_{x \in \Delta_m} \tilde{\Phi}^{\min}(x) = \max_x \min_{y \in \Delta_n} x^T M y = \tilde{v}^-$$

N.B.. $\tilde{\Phi}$ is LINEAR if $y \notin \Delta_n$ is POLYHEDRON \Rightarrow

$\min_{y \in \Delta_n} x^T M y$ is attained at vertex, i.e. at $(0, \dots, 0, 1, 0, \dots, 0)$

i.e. pure strategy. \Rightarrow

$$\tilde{\Phi}^{\min}(x^*) = \max_{x \in \Delta_m} \underbrace{\min_{j=1, \dots, n} (x^T M)_j}_{\tilde{\Phi}^{\min}(x)}$$

Similarly y^* is S.S. \Leftrightarrow

$$\tilde{\Phi}^{\max}(y^*) = \min_y \max_{i=1, \dots, m} (M y)_i$$

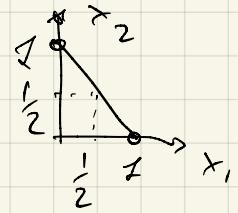
$$\left(\begin{array}{c} x^T M y \\ x = (0, \dots, 0, 1, 0, \dots, 0) \end{array} \right)$$

Ex. 1 Head & tail game or EVEN-ODDS with 2 fingers.

$$M = \begin{pmatrix} 1 & -1 \\ -1 & 1 \end{pmatrix} \quad (x_1, x_2) M = (\underbrace{x_1 - x_2}, \underbrace{-x_1 + x_2})$$

$$\tilde{\Phi}^{\min}(x) = \min \left\{ \underbrace{x_1 - x_2}, \underbrace{-x_1 + x_2} \right\} = -|x_1 - x_2|$$

$$\max_{x \in \Delta_2} \tilde{\Phi}^{\min}(x) = \max_{\Delta_2} \underbrace{(-|x_1 - x_2|)}_{\leq 0} \stackrel{!}{=} 0 \text{ iff } x_1 = x_2$$



$x^* = (\frac{1}{2}, \frac{1}{2})$ is the only sec. strat.

$$\tilde{\Phi}^{\max}(y) = |y_1 - y_2| \geq 0 \quad \not{=} 0 \iff y_1 = y_2 = \frac{1}{2}$$

So $y^* = (\frac{1}{2}, \frac{1}{2})$ is SEC. STRAT.

Conclusion: $((\frac{1}{2}, \frac{1}{2}), (\frac{1}{2}, \frac{1}{2}))$ is the UNIQUE saddle pt. in mixed strategies & $\tilde{v} = 0$. □

Example ROCK - PAPER - SCISSOR.

	R	P	S	$\min \rightarrow$
R	0	-1	1	-1
P	1	0	-1	-1
S	-1	1	0	-1

$\max \downarrow$

1 1 1

$\hookrightarrow v^+ = 1$

$\min \swarrow$

$v^- = -1$

$v \neq 0$, i.e. pure strategies.

Look for \tilde{v} in mixed strategies in $\Delta_3 \times \Delta_3$

$$\left\{ \begin{array}{l} \tilde{\Phi}^{\min}(x^*) = \max_{x \in \Delta_3} \tilde{\Phi}^{\min}(x) = \max_x \min_{j=1, \dots, 3} (x^T M)_j \\ \tilde{\Phi}^{\max}(y^*) = \min_y \tilde{\Phi}^{\max}(y) = \min_{y \in \Delta_3} \max_{i=1, \dots, 3} (M y)_i \end{array} \right.$$

$$x^T M = \begin{pmatrix} x_1 & x_2 & x_3 \end{pmatrix} \begin{pmatrix} 0 & -1 & 1 \\ 1 & 0 & -1 \\ -1 & 1 & 0 \end{pmatrix} = (x_2 - x_3, -x_1 + x_3, x_1 - x_2)$$

$$\tilde{\Phi}^{\min}(x) = \min_{\Delta_3} \{x_2 - x_3, x_3 - x_1, x_1 - x_2\}$$

Note that $\tilde{\Phi}^{\min}(x) \leq 0$ if x because, if not,

$$\begin{cases} x_2 > x_3 \\ x_3 > x_1 \\ x_1 > x_2 \end{cases} \quad \text{ACROSS, B.L.E! Then search } x^* : \quad \tilde{\Phi}^{\min}(x^*) = 0 \quad \text{I get: } x_2 = x_3 = x_1$$

$$x_2 + x_3 + x_1 = 1 \quad \Rightarrow \quad x_1 = \frac{1}{3} = x_2 = x_3 \quad x^* = \left(\frac{1}{3}, \frac{1}{3}, \frac{1}{3}\right) \text{ is SEC. STRAT. (UNIQUE),}$$

H.W. Similarly $\tilde{\gamma}^*(\frac{1}{3}, \frac{1}{3}, \frac{1}{3})$ is the only S.S. for γ .

CONCLUSION: (x^*, y^*) is the only SADDLE point of RPS game in MIXED STRATEGIES.

HW Ex [Banach Thm. 2.11 p.56-58]

Every 2×2 game $M = \begin{pmatrix} \phi_{11} & \phi_{12} \\ \phi_{21} & \phi_{22} \end{pmatrix}$ has a saddle point

EITHER in PURE STRAT.s or in MIXED STR. COMPUTABLE
Explicitly by solving 2 linear equations.

Value in mixed strategies for general 2-p 0-sum games. (just main ideas, details in [Bresser]).

$\tilde{\Phi}: A \times B \rightarrow \mathbb{R}$ cont., $A, B \subseteq \mathbb{R}^k$ compact

$\mathbb{P}(A), \mathbb{P}(B)$ subsets of vector space of Radon measures.

Lemma. $\mathbb{P}(A), \mathbb{P}(B)$ are convex.

Pf. $\lambda \in [0, 1]$, $\bar{\mu}_1, \bar{\mu}_2 \in \mathbb{P}(A)$, $\lambda \bar{\mu}_1 + (1-\lambda) \bar{\mu}_2 \geq 0$

$$(\lambda \bar{\mu}_1 + (1-\lambda) \bar{\mu}_2)(A) = \lambda \bar{\mu}_1(A) + (1-\lambda) \bar{\mu}_2(A) = 1. \quad \square$$

Topology on $\mathbb{P}(A) \neq \mathbb{P}(B)$ that makes them metrizable & compact : weak \star topology on $\mathbb{P}(A)$ as a dual of $C(A)$ with uniform convergence.

Def. $(\mu_k)_{k \in \mathbb{N}}$ in $\mathbb{P}(A)$ conv. weak \star to $\mu \in \mathbb{P}(A)$ if.

$$\forall f \in C(A) \quad \lim_{k \rightarrow \infty} \int_A f d\mu_k = \int_A f d\mu.$$

& we write $\mu_k \xrightarrow{\star} \mu$.

Rmk. $\tilde{\Phi} : \mathbb{P}(A) \times \mathbb{P}(B) \rightarrow \mathbb{R}$ ($\tilde{\Phi}$ is cont.) is cont. v.n.f.

the product weak \star convergence. i.e., $\mu_k \xrightarrow{\star} \mu, \nu_k \xrightarrow{\star} \nu$

$$\lim_{k \rightarrow \infty} \tilde{\Phi}(\mu_k, \nu_k) = \lim_{k \rightarrow \infty} \iint_{A \times B} \tilde{\Phi}(a, b) d\mu_k(a) d\nu_k(b) =$$

$$= \iint_{A \times B} \tilde{\Phi}(a, b) d\mu(a) d\nu(b) = \tilde{\Phi}(\mu, \nu), \quad (\text{No PROOF})$$

Thm A compact $\Rightarrow P(A)$ with \cong topology is METRIZABLE & SEQUENTIALLY COMPACT., i.e., $\{\mu_{k_n}\}_{n \in \mathbb{N}}$ in $P(A)$, $\exists \mu_{k_n}, k_n \rightarrow \infty : \mu_{k_n} \xrightarrow{n \rightarrow \infty} \mu \in P(A).$

Pf No

Cor (of V.Neumann Min-Max Thm). $A, B \subseteq \mathbb{R}^k$ compact, $\Phi \subseteq C(A \times B)$ \Rightarrow the game has a value in mixed strategies.

Pf. Game $(P(A), P(B), \tilde{\Phi})$ satis. ass. of V.Ne.Thm.

Remaining only: $\tilde{\Phi}$ is CONCAVE-CONVEX.

$$\begin{aligned} \tilde{\Phi}(\lambda \mu + (1-\lambda) \bar{\mu}, \nu) &= \iint_{A \times B} \tilde{\Phi}(a, b) (\lambda d\mu(a) + (1-\lambda) d\bar{\mu}(a)) d\nu(b) = \\ &= \lambda \iint_{A \times B} \tilde{\Phi}(a, b) d\mu(a) d\nu(b) + (1-\lambda) \iint_{A \times B} \tilde{\Phi}(a, b) d\bar{\mu}(a) d\nu(b) = \\ &= \lambda \tilde{\Phi}(\mu, \nu) + (1-\lambda) \tilde{\Phi}(\bar{\mu}, \nu) \quad \Rightarrow \mu \mapsto \tilde{\Phi}(\mu, \nu) \text{ is LINEAR} \end{aligned}$$

\Rightarrow it is CONCAVE.

Similarly $\nu \mapsto \tilde{\Phi}(\mu, \nu)$ is LINEAR \Rightarrow CONVEX.

V.Ne.Thm $\Rightarrow (P(A), P(B), \tilde{\Phi})$ has a solnle. \square



NON-ZERO SUM GAMES.

For simplicity just 2 players

Given A, B sets of decisions or strategies.

$\tilde{\Phi}^A : A \times B \rightarrow \mathbb{R}$ payoff. of 1st player, want. to fix it

$\tilde{\Phi}^B : A \times B \rightarrow \mathbb{R}$. $n \leftarrow 2^{\omega} n$, $n = 2$

- ZERO SUM , $\Phi^A + \Phi^B = 0$ $\forall \alpha \rightarrow \Phi^B = -\Phi^A$
- if $\Phi^A - \Phi^B = 0$ leads to OPTIMIZ. PB. in $A \times B$