

# LECTURE 15, April 27, 23

GAME THEORY:  $\Phi \in C(A \times B)$ ,  $A, B$  compact metric spaces.

MARGINAL FNS:  $\Phi^{\max}(b) := \max_a \Phi(a, b)$ ,  $\Phi^{\min}(a) := \min_b \Phi(a, b)$

Best Response:  $R^A(b) := \arg \max_a \Phi(a, b)$ ,  $R^B(a) = \arg \min_b \Phi(a, b)$

UPPER VALUE:  $v^+ := \min_b \Phi^{\max}(b) = \min_b \max_a \Phi(a, b)$

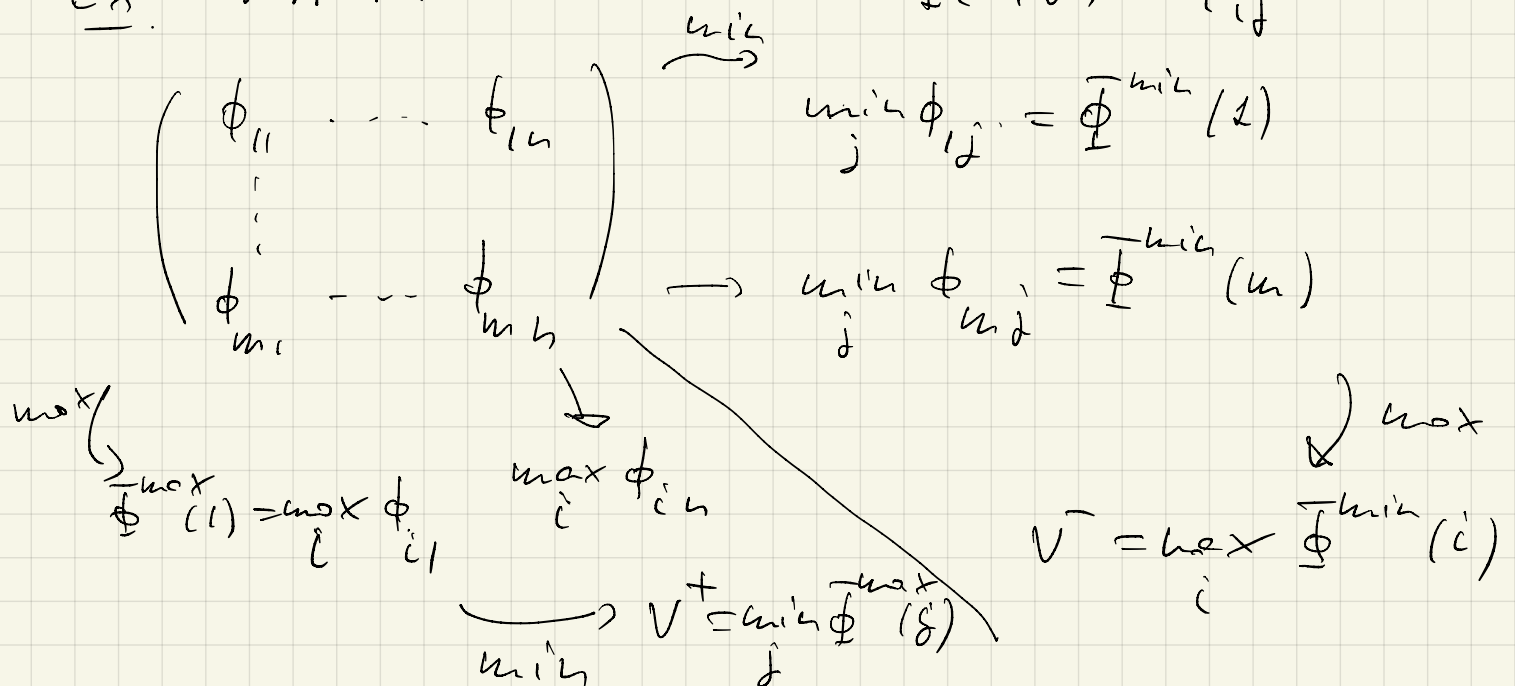
LOWER VALUE:  $v^- := \max_a \Phi^{\min}(a) = \max_a \min_b \Phi(a, b)$

Prop.  $v^- \leq v^+$ . If  $v^- = v^+$  GAME HAS A VALUE.

Examples of  $v^- \neq v^+$ .

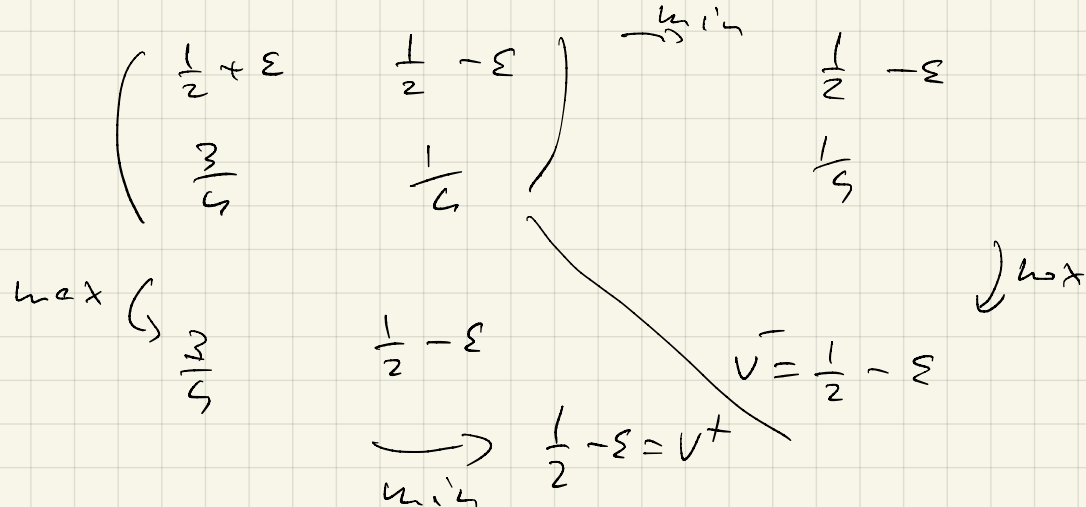
Ex. MATRIX GAMES

$$\Phi(i, j) = \phi_{ij}$$



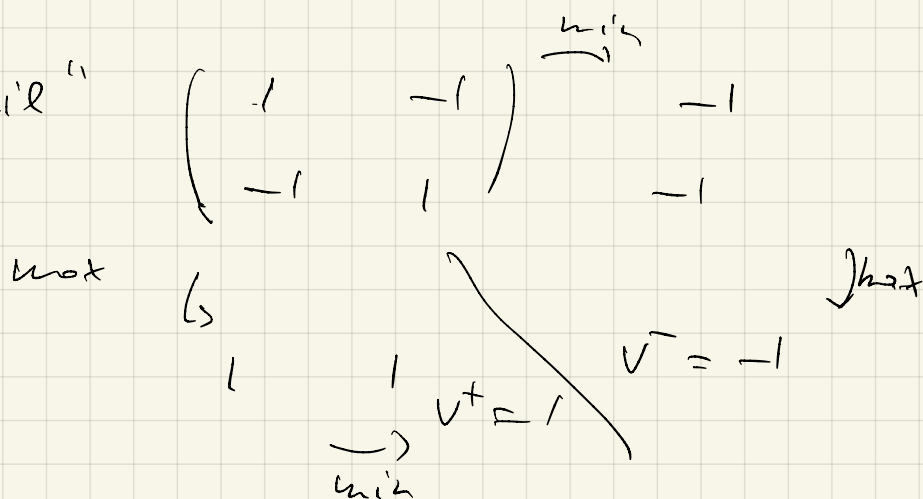
Ex. 2 "Cake"

$$0 < \varepsilon < \frac{1}{4}$$



$\Rightarrow V = \frac{1}{2} - \varepsilon$  is THE VALUE of the game.

Ex. 3 "Head & Tail"



$$V^- = -1 < V^+ = 1$$

The VALUE does NOT exist!

Def. A SADDLE POINT of the game is  $(a^*, b^*) \in A \times B$ :

$$\forall a \in A \quad \Phi(a, b^*) \leq \Phi(a^*, b^*) \leq \Phi(a^*, b) \quad \forall b \in B$$

Ex.  $A = B = [-1, 1] \quad \Phi(a, b) = b^2 - a^2$



$b^* = a^* = 0$  is a SADDLE

Rank:  $(a^*, b^*)$  is a saddle  $\Leftrightarrow$

$$(S) \quad \max_a \Phi(a, b^*) \leq \Phi(a^*, b^*) \leq \min_b \Phi(a^*, b)$$

$a^* \in R^A(b^*) \qquad b^* \in R^B(a^*)$

$$(S) \Leftrightarrow (S') \quad \max_a \Phi(a, b^*) = \Phi(a^*, b^*) = \min_b \Phi(a^*, b)$$

Rank. Suppose  $R^A$  &  $R^B$  are functions (single-valued)

$R^A: B \rightarrow A$ ,  $R^B: A \rightarrow B$ .  $(a^*, b^*)$  saddle  $\Rightarrow$

$a^*$  is a fixed point of  $R^A \circ R^B: A \rightarrow A$  because

$$R^A \circ R^B(a^*) = R^A(b^*) = a^*$$

&  $b^*$  is a fixed point of  $R^B \circ R^A: B \rightarrow B$ .  $\square$

Def. SECURITY STRATEGIES:  $v^+ = \min_b \Phi^{\max}(b)$

is  $b^*$ :  $v^+ = \Phi^{\max}(b^*)$ , i.e.,  $b^* \in \arg \min_b \Phi^{\max}(b)$   
for 2<sup>nd</sup> player.

$a^*$  is S.S. for 1<sup>st</sup> player if  $v^- = \Phi^{\min}(a^*)$

i.e.,  $a^* \in \arg \max_a \Phi^{\min}(a)$

Thm. The game has a value  $\Leftrightarrow \exists$  a saddle point.

Pf " $\Leftarrow$ " ASS.:  $(a^*, b^*)$  saddle pt. Goal:  $v^+ \leq v^-$

$$v^- = \max_a \min_b \Phi(a, b) \geq \min_b \Phi(a^*, b) = \max_a \Phi(a, b^*)$$

$\geq \min_b \max_a \Phi(a, b) = v^+$   $\square$

Remark  $a^*$  is a SEC. STR. for 1<sup>st</sup> player.  $b^*$  is S.S. for 2<sup>nd</sup>

" $\Rightarrow$ " ASS  $v^+ = v^-$ . Take  $a^* \in$  SEC. STR. for A.

$$v^- = \Phi^{\min}(a^*) = \min_b \Phi(a^*, b)$$

Take  $b^* \in$  SEC. STR. for B  $\therefore v^+ = \Phi^{\max}(b^*) = \max_a \Phi(a, b^*)$

$$\forall a \in A \quad \Phi(a, b^*) \leq \max_a \Phi(a, b^*) = v^+ \stackrel{\text{ASS}}{=} v^- = \min_b \Phi(a^*, b) \\ \leq \Phi(a^*, b) \quad \forall b$$

$$\text{For } a = a^*, b = b^* \Rightarrow \Phi(a^*, b^*) = v^+ = v^- \leq \Phi(a^*, b^*)$$

$$\Rightarrow \text{"} \leq \text{" or " = " } \Rightarrow \Phi(a^*, b^*) = v \quad \# \quad (S) \text{ holds,}$$

$$\Rightarrow (a^*, b^*) \text{ is a SADDLE PT.} \quad \#$$

Corollary If game has a value  $\Rightarrow$

(i)  $(a^*, b^*)$  is a saddle  $\Leftrightarrow a^*$  is SEC. STR. for 1<sup>st</sup>  $b^*$  is S.S. for 2<sup>nd</sup>

(ii) (EXCHANGEABILITY) : If  $(\bar{a}, \bar{b})$  is also a saddle.

$$\Rightarrow (a^*, \bar{b}), (\bar{a}, b^*) \text{ are saddles.}$$

Pf (ii) See pf. of thm. (ii) from (i).  $\#$

THE MINIMAX THEOREM of Von Neumann (1924)

Thm:  $A, B \subseteq$  vector spaces, COMPACT & CONVEX,  $\Phi \in C(A \times B)$

$\Phi$  CONCAVE-CONVEX, i.e.,

$$\left\{ \begin{array}{l} \forall b, \quad a \mapsto \Phi(a, b) \text{ is CONCAVE} \\ \forall a, \quad b \mapsto \Phi(a, b) \text{ is CONVEX} \end{array} \right.$$

$\Rightarrow V^+ = V^-$ , i.e.,  $(A, B, \Phi)$  has a value  $\Phi$  at least one SADDLE POINT.

PROOF: will be done if  $A \subseteq \mathbb{R}^m, B \subseteq \mathbb{R}^n$ .

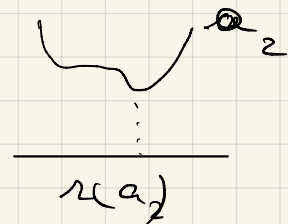
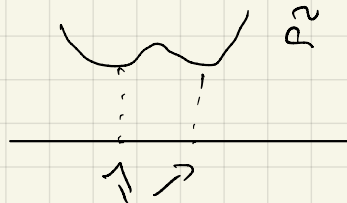
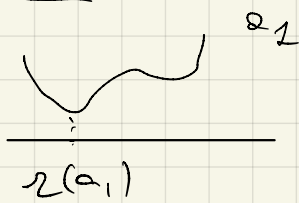
RMK. Supp.  $\Phi(a, \cdot)$  STRICTLY CONVEX  $\Rightarrow \exists$  UNIQUE

$$r(a) \in B : \Phi(a, r(a)) = \min_b \Phi(a, b)$$

$$\Rightarrow R^B(a) = \{r(a)\}.$$

Lemma. If  $\Phi(a, \cdot)$  is strictly convex  $\forall a \in A, \Phi \in C$ ,  $A, B$  compact, convex  $\Rightarrow r: A \rightarrow B$  is continuous.

RMK. Not true without convexity.



here  $r$  jumps to right at  $\bar{a}$ .  $R^B(\bar{a})$

Pf Lemma.  $\bar{a} \in A, a_n \rightarrow \bar{a}$  goal:  $r(a_n) \rightarrow r(\bar{a})$

Extract  $a_{n_k} : r(a_{n_k}) \rightarrow \bar{b} \in B$  ( $B$  compact.)

$$\Rightarrow \Phi(a_{n_k}, r(a_{n_k})) \rightarrow \Phi(\bar{a}, \bar{b}) \quad k \rightarrow \infty$$

$$\forall b \quad \left( \leq \Phi(a_{n_k}, b) \right) \rightarrow \Phi(\bar{a}, b) \quad "$$

$$\Rightarrow \Phi(\bar{a}, \bar{b}) \geq \Phi(\bar{a}, \bar{b}) \quad \forall \bar{b} \in B$$

$$\Rightarrow \bar{b} = r(\bar{a}) \quad \Rightarrow \quad r(a_{n_k}) \xrightarrow{b \rightarrow \infty} r(\bar{a})$$

By the arbitrariness of  $a_{n_k} \Rightarrow r(a_n) \xrightarrow{n \rightarrow \infty} r(\bar{a})$  □

Pf of v.N. Theorem: Step 1: Supp.  $\forall a \in A$   $b \mapsto \Phi(a, b)$  is strictly convex. Then  $\forall a \exists r(a)$  CONT.

$$\Phi(a, r(a)) = \min_b \Phi(a, b)$$

Step 2: Take  $a^+$  SEC. STRAT. for 1<sup>st</sup> pl. , i.e. ,

$$v^- = \max_a \Phi^{\min}(a) = \Phi^{\min}(a^+) \quad b^* := r(a^+)$$

Goal:  $(a^+, b^*)$  is a saddle pt.

$$\text{Note: } \Phi(a^+, b^*) = \min_b \Phi(a^+, b) \leq \Phi(a^+, b) \quad \forall b$$

$\uparrow r(a^+)$

Remains the goal.  $\Phi(a^+, b^*) \geq \Phi(a, b^*) \quad \forall a$ .

Step 3: Idea: approximate  $a^+$  with  $a_d := \lambda a + (1-\lambda)a^+ \xrightarrow{\lambda \rightarrow 1} a^+$  (if  $\lambda \rightarrow 0$ )

fix  $a \in A$ ,  $\lambda \in [0, 1]$ ,  $\mu = 1 - \lambda$

$$\Phi(a^+, b^*) \geq \Phi^{\min}(a^+) \geq \Phi^{\min}(a_d) =$$

$\underbrace{a^+ \text{ is SEC. STRAT.}}_{\text{circled}}$

$$= \Phi(a_d, r(a_d)) \geq \lambda \Phi(a, r(a_d)) + \mu \Phi(a^+, r(a_d))$$

$\Phi$  conc. in  $a$ .

$$\geq \lambda \Phi(a, r(a_d)) + (1-\lambda) \Phi^{\min}(a^+)$$

$$\Rightarrow \cancel{\lambda} \Phi^{\min}(a^+) \geq \cancel{\lambda} \Phi(a, r(a_d)) \quad \downarrow \lambda \rightarrow 1$$

$$\Rightarrow \Phi^{\min}(a^*) \geq \Phi(a, z(a^*)) \quad \text{USE LEMMA 1}$$

$\stackrel{z}{=} b^*$

$$\Rightarrow \Phi^{\min}(a^*) \geq \Phi(a, b^*) \quad \forall a$$

which is the goal  $\square$  st. 2

$\Phi(a^*, b^*) \geq$

Step 4. Remove the strict convexity of  $b \mapsto \Phi(a, b)$ .

For simplicity here  $B \subseteq \mathbb{R}^k$ . Fix  $\varepsilon > 0$ :

$$\Phi_\varepsilon(a, b) = \Phi(a, b) + \varepsilon |b|^2 \quad \text{is STRICTLY CONVEX in } b \quad \forall a.$$

st. 2-3  $\Rightarrow \Phi_\varepsilon$  has a saddle point  $(a_\varepsilon, b_\varepsilon)$ , i.e.

$$\forall a \quad \Phi_\varepsilon(a, b_\varepsilon) \leq \Phi_\varepsilon(a_\varepsilon, b_\varepsilon) \leq \Phi_\varepsilon(a_\varepsilon, b) \quad \forall b.$$

By compactness of  $A, B$ , extract  $\varepsilon_n \searrow 0$ :

$$a_{\varepsilon_n} \rightarrow a^* \in A, \quad b_{\varepsilon_n} \rightarrow b^* \in B.$$

$$\begin{aligned} \Phi_{\varepsilon_n}(a, b_{\varepsilon_n}) &\leq v_{\varepsilon_n} = \Phi_{\varepsilon_n}(a_{\varepsilon_n}, b_{\varepsilon_n}) = \Phi(a_{\varepsilon_n}, b_{\varepsilon_n}) + \varepsilon_n |b_{\varepsilon_n}|^2 \\ &\leq \Phi(a_{\varepsilon_n}, b) + \varepsilon_n |b|^2 \quad \forall b \end{aligned}$$

$$\text{let } \varepsilon_n \rightarrow 0 \quad \Phi(a, b^*) \leq \Phi(a^*, b^*) \leq \Phi(a^*, b) \quad \forall b$$

$\Rightarrow (a^*, b^*)$  is a SADDLE.  $\square$

Examples. 1:  $\Phi(a, b) = \varphi_1(a) - \varphi_2(b)$

$\varphi_1, \varphi_2$  CONCAVE in  $[1, 1]^m \Rightarrow$  thm. applies.

2 IMPORTANT:  $\Phi(a, b) = a^T M b$   $M \in \mathcal{M}_{m \times n}$

$a \in A \subseteq \mathbb{R}^m, b \in B \subseteq \mathbb{R}^n$ .  $\Phi$  is bilinear  $\Rightarrow$   
 $\uparrow$   $\xrightarrow{\quad}$   
conv.  $\Phi$  convex cont.  $\Phi$  linear in  $a$   $\Phi$  l.s

$\Rightarrow \Phi$  CONC. CONVEX THM. is OK.

3. MATRIX GAMES  $A = \{1, \dots, m\}, B = \{1, \dots, n\}$ .  
are NOT CONVEX, V.N.T. does NOT apply  $\Phi$  in  
fact we know examples without value.

MIXED STRATEGIES. Idea: choose in a stochastic  
instead of deterministic way.

Def. A mixed strategy of 1<sup>st</sup> player is a  $\mu \in \mathcal{P}(A) :=$   
 $= \{ \text{probability measures on } A \}$ . and for 2<sup>nd</sup> player it is  
 $\nu \in \mathcal{P}(B) = \{ \dots \text{ on } B \}$ .

Ex:  $\delta_{\bar{a}}$  = Dirac measure concentrated in  $\bar{a} \in A$ . c. e.

$\forall S \subseteq A$  Borel  $\delta_{\bar{a}}(S) = \begin{cases} 1 & \text{if } \bar{a} \in S \\ 0 & \text{if } \bar{a} \notin S \end{cases}$

$\mathcal{P}(A) \cong$  "copy of  $A$ ".  $A$  = pure strategies

Def  $\tilde{\Phi}(\mu, \nu) := \iint_{A \times B} \Phi(a, b) d\mu(a) d\nu(b)$

$\tilde{\Phi}: \mathcal{P}(A) \times \mathcal{P}(B) \rightarrow \mathbb{R}$ .



N.B.  $\tilde{\Phi}(\delta_{\bar{a}}, \delta_{\bar{b}}) = \iint_{A \times B} \Phi(a, b) d\delta_{\bar{a}}(a) d\delta_{\bar{b}}(b) =$   
 $= \Phi(\bar{a}, \bar{b})$

$\int_A f(a) d\delta_{\bar{a}}(a) = f(\bar{a}).$

$\Rightarrow \tilde{\Phi}$  "EXTENDS"  $\Phi$  from  $A \times B$  to  $P(A) \times P(B)$

Def. If  $\exists$  value of the game  $(P(A), P(B), \Phi)$  & a saddle pt., they are called value & saddle of  $(A, B, \Phi)$  in mixed strategies.